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ALGAE I\*

A Compiler for the IBM 704†

by

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†This version is for use with a machine having: 8.192 words of core storage, 4 logical drums, 7 tape units, CAD instruction. A peripheral tape-to-printer or simulator is required.

## 1. Introduction

Discussions involving the subject of defining problems for interpretation and coding by known automatic-coding systems generally suggest that the techniques for stating the control (or logic) of the problem are frequently difficult to understand and difficult to use. It seems that the difficulty is one of discovering a suitable language with which to define problem control. In programming a problem for hand coding, the familiar flow diagram has been successfully used (when needed) to define the control of the problem. Such flow diagrams (as we draw them) cannot be presented directly to present-day computers. It is the purpose of this paper to propose a flow diagram representation using simple algebraic language which can be directly entered into the computer and to describe a compiler which accepts problems coded in this form.

The first part of this description concerns itself with the development of an idea which was earlier presented in a paper entitled "Algebraic Formulation of Flow Diagrams". The second section describes a compiler for the IBM 704 (called ALGAE I) which incorporates most of the techniques described in the first part. ALGAE I makes use of FORTRAN I to form the actual 704 code.

ALGAE I was written as a means to permit programmers to use and evaluate this idea of program control. As a result, any comments, criticisms, etc., from users of the system are not only welcome, but eagerly solicited.

## 2. General Description of the Algol Language.

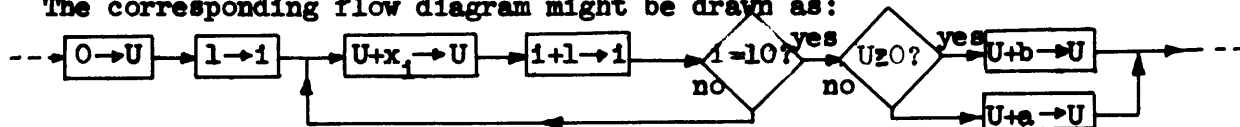
### 2.1 Basic Concepts.

Assume we have a flow diagram such that it could be used to hand-code a problem for any stored-program computer. Suppose we remove from the boxes all equations, statements of input-output tasks and other statements not directly related to the control and logic of the problem, and list them, with identification, elsewhere as reference material. We do not include in this list statements and questions pertaining to loops, numerical conditions, or switch and trigger conditions. The items remaining in the flow-diagram would form a statement or "part-picture" of the control for the problem. We will attempt to translate this control statement into a statement which can be easily written and entered into a computing machine.

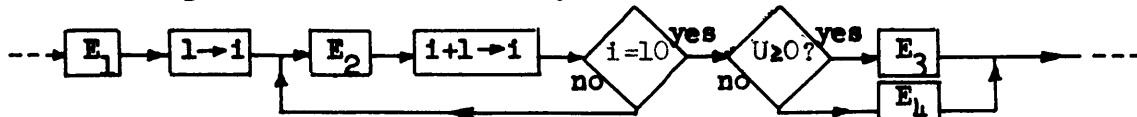
In order to illustrate the above discussion and extend the idea further, let us consider an elementary example. Suppose in a problem we wished to compute

$$U = \sum_{i=1}^{10} x_i + \begin{cases} a & (\text{if } \Sigma < 0) \\ \text{or} \\ b & (\text{if } \Sigma \geq 0) \end{cases}$$

The corresponding flow diagram might be drawn as:



Removing the items not directly related to control would leave:



If a distinction is now made between loop ranges, such as the  $i$ -loop in the above example, and conditions for execution of equations, such as the  $U$  test above, we introduce the symbols

$$I_1 : i = 1(1)10 \quad (\text{meaning } i \text{ takes successively the values } 1 \text{ through } 10 \text{ in increments of } 1)$$

and  $C_1 : U \geq 0$ .

Our example would then be completely defined as follows:

Example 1

<u>Control</u>	<u>Equations</u>
$\dots + E_1 + I_1 E_2 + C_1(E_3, E_4) + \dots$	$E_1 : U = 0$
$I_1 : i = 1(1)10$	$E_2 : U = U + x_i$
$C_1 : U \geq 0$	$E_3 : U = U + b$
	$E_4 : U = U + a$

where  $C_1$  here is understood to mean "execute  $E_3$  if  $C_1$  is satisfied and execute  $E_4$  if  $C_1$  is not satisfied." If we were to define  $C_1^*$  as the negation of the condition  $C_1$ , the last term in the control statement could be written as

$$\dots + C_1 E_3 + C_1^* E_4 + \dots$$

provided it is known that when  $E_3$  is executed  $E_4$  would not also be executed (as would be the case if  $\sum x_i = 1$  and  $b = -2$ ). Clearly the  $C^*$  convention is not essential since the condition  $C_2 : U < 0$  could be used instead. The meaning of  $I_1 E_2$  is evident and will be discussed more generally below.

## 2.2 Definition of Symbols.

Let us now define more completely the set of control statement symbols. These definitions make no restrictions on the characters with regard to their use in equation writing since control statements are assumed to be handled separately from equations.

C (assumed to have a subscript) represents a single condition for a two-way branch. Several conventions, defined in section 2.5, have been developed which use logical combinations of C's. C's are used for stating all conditions, except those inherent in range statements such as end-of-loop tests. C-type conditions include tests of numerical, logical, console switch, and trigger conditions. A more generalized, n-way branch is under development. Once a C has been defined, it may be used repeatedly throughout the problem.

E (assumed to have a subscript) represents a single equation, a continuous and closed set of equations (whenever one is done, all are done), an input-output task, program stop, or any other

statement or closed set of statements not directly related to problem control. An  $E_i$  may also be used more than once.

$G$  and  $H^\dagger$  (subscripted). This convention is designed to facilitate transfer of control. The symbol  $G$  (with a subscript) denotes a transfer to  $H$  (with the same subscript). There may be a many-to-one correspondence between  $G$  and  $H$ .

$I, J, K, L, M, N$  (subscripted) represent single range statements or loop definitions.  $I_1$  in Example 1 represents the range statement  $i = 1(1)10$ . More generally, the expression  $I_j : a = b(c)d$  means that  $I_j$  defines a loop in which the subscript  $a$  varies from  $b$  to  $d$  in steps of  $c$ . These may be used repeatedly.

$S$  (subscripted) represents a control statement. It allows the program control to be defined by many sub-control statements which are in turn connected by a single master control statement. An  $S_i$ , once defined, may be used only one time. This partitioning of the code into smaller units of control tends to clarify the various phases of the problem.

$T$  (subscripted) is basically an  $S$  which may be used repeatedly within a problem. It is the Algebraic representation of a closed subroutine. Both the  $S_i$  and  $T_i$  can appear in the same positions as an  $E_i$ .

The plus sign, (+), has two meanings. The first is to indicate a direct flow of control, as, for example,  $E_1 + E_2$ , which means "execute  $E_1$ , then  $E_2$ ."  $E_1$  and  $E_2$  are terms of the control statement  $E_1 + E_2$ . A second meaning, when used with  $C$ 's, is discussed in section 2.5.

The parentheses, ( ), are used for the phrasing or grouping of terms, for indicating ranges, and for special purposes to be described. The expression  $(E_1 + E_2)$  is a single term.

The comma is used in special conventions, as described below.

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<sup>†</sup>This set of symbols has not been shown to be essential to the system. It serves more as a convenience to the coder and may introduce some undesirable flexibilities unless used with caution.

### 2.3 The Basic Forms.

The symbols defined above may appear in seven basic forms. Section 2.4 sets forth rules to be used in combining these basic forms to symbolize more complex logical conditions.

Let us define the quantity X as any single subscripted E, S, or T. The basic forms may be written as:

1. X            The single quantity. In this case X may also be a subscripted G or H.
2.  $I_j X$         Execute X loopwise for the range of the subscript indicated in the definition of  $I_j$ .
3.  $C_1 X$         Execute X only if the condition  $C_1$  is satisfied. In this form, X may be the subscripted transfer symbol  $G_1$ , resulting in a conditional transfer.
4.  $X C_1$         Execute X, test  $C_1$ ; if  $C_1$  is not satisfied, execute X again, test  $C_1$  and continue iterating until  $C_1$  is satisfied.
5.  $C_j(X_a, X_b)$     If  $C_j$  is satisfied, do  $X_a$  and skip  $X_b$ . If  $C_j$  is not satisfied, skip  $X_a$  and execute  $X_b$ .  $X_a$  and/or  $X_b$  may be the subscripted symbol  $G_1$  in this form.
6.  $(X_a, X_b) C_j$     Execute  $X_b$ , test  $C_j$ . If  $C_j$  is not satisfied, do  $X_a$  and  $X_b$ , test  $C_j$  and iterate until  $C_j$  is satisfied. In a control sense, this is an iteration loop in which the statement,  $X_a$ , is executed in every iteration except the first.
7.  $I_j(X) C_k$       After each cycle of the  $I_j X$  loop, test  $C_k$  and if satisfied, leave the loop and calculate the next term in the control statement. If  $C_k$  is not satisfied, do the next cycle, test  $C_k$  again, etc. Exit from the  $I_j X$  loop is made as in form 2 unless  $C_k$  is satisfied earlier.

In each of the above forms involving  $I_j$ , any other range statement could be used ( $N_k$ ,  $L_1$ , etc.). Note that all characters must be subscripted.

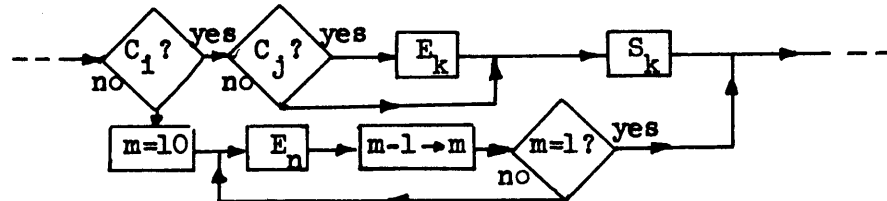
2.4 Combinations.

More complex operations may be expressed by combinations of the above forms in which the following rules are applied.

1. Any basic form may be combined with another by using the "continuation symbol" or plus sign (+). Thus,  $E_1 + S_k$  means "execute the equations in  $E_1$  as defined, then perform the sequence of operations defined by  $S_k$ ."
2. Any basic form may be placed within another basic form (or within itself) in the position occupied by the symbol X of region 2.3. Thus  $I_k(C_j E_1)$ , form 3 within form 2, means "for each cycle of the  $I_k(\dots)$  loop, calculate  $E_1$  only if  $C_j$  is satisfied; otherwise start the next cycle." Or,  $I_k(J_1 E_k)$ , form 2 within itself, means "for each value of the subscript whose range is defined by  $I_k$ , execute the loop defined by  $J_1 E_k$ ." This is the familiar loop within a loop.

Note that the basic form which replaces the X is set off by parentheses to indicate order of performance. Parentheses are not required in forms 5,6, or 7 since the form itself sets off the X's by parentheses and commas.

3. Any meaningful combination created by the above two rules may also be inserted within a basic form in place of an X. This "second-order" combination may be illustrated by the term  $C_1(C_j E_k + S_k, I_k E_n)$ , forms 3 and 1 combined using rule 1 and placed, with form 2, within form 5. The meaning of this term can best be illustrated by the flow diagram



where  $I_k$  is defined, for convenience of illustration, as  $m = 10(-1)2$  and  $C_1?$  asks the question, "Is the  $C_1$  condition satisfied?"

Additional compounding may be achieved by repeated use of the above rules.

## 2.5 C-Statements.

Further substitution may be made in the combinations and basic forms by replacing a single C condition with a C-Statement. Since the C's are defined as binary conditions, a logically arranged set of C's can have the effect of a single C. For this, we use the second definition of the symbol (+) and say that  $C_1 + C_j$  means  $C_1$  or  $C_j$ . Similarly, let us define  $C_1 C_j$  as  $C_1$  and  $C_j$ . Substituting an expression consisting only of C's into a basic form or combination allows us to express conditions such as  $(C_1 + C_2 C_3)E_1$ , meaning "if  $C_1$  is satisfied or if  $C_2$  and  $C_3$  are both satisfied, execute  $E_1$ ; otherwise skip  $E_1$  and continue to the next term."

C-statements may also contain negated conditions by using the negation symbol (\*). Thus  $(C_1 + C_2 C_3)^*$  would be the negation of the total situation defined in the above example. More complex statements of this type may be defined by the basic rules of Boolean Algebra.

When substituting a C-statement into a form or combination, parentheses are needed to encompass the statement only if the resulting terms are OR-ed at some point. Thus  $C_3 C_4 + (C_2 + C_5)$  should be enclosed by parentheses as  $(C_3 C_4 + (C_2 + C_5))$ , but the term  $C_3 C_4 (C_2 + C_5)$  need not be.

To illustrate, re-define example 1 as:

$$U = \sum_{i=1}^{10} x_i + \left\{ \begin{array}{l} a \text{ (if } \Sigma > 0, \Sigma \leq 3 \text{ and } x_1 \geq 4) \\ b \text{ (if } \Sigma > 3 \text{ or } x_1 < 4) \\ 0 \text{ otherwise.} \end{array} \right\}$$

We could completely define this as follows:

Example 1a

<u>Control</u>	<u>Equations</u>
$\dots + E_1 + I_1 E_2 + C_1 C_2 C_3^* (E_4, (C_2^* + C_3) E_3) + \dots$	$E_1 : U = 0$
$I_1 : i = 1(1)10$	$E_2 : U = U + x_1$
$C_1 : \Sigma > 0$	$E_3 : U = U + b$
$C_2 : \Sigma \leq 3$	$E_4 : U = U + a$
$C_3 : x_1 < 4$	



## 2.6 Sample problem.

Let us now illustrate some of the preceding ideas with another example. Assume we wish to solve the Laplace equation for a 10 x 10 mesh with  $x = 1$  on the boundaries and  $x = 0$  in the interior. Then:

### Example 2.

<u>Control</u>	
$S_1 : I_1(J_1((C_1 + C_2 + (C_3 C_4)^*)(E_1, E_2)))$	
$S_2 : E_3 + I_2(J_2 E_4)$	
$S_3 : S_1 + S_2 C_5 + E_5$	
$I_1 : i = 1(1)10$	$C_1 : i = 1$
$I_2 : i = 2(1)9$	$C_2 : i = 10$
$J_1 : j = 1(1)10$	$C_3 : j \geq 2$
$J_2 : j = 2(1)9$	$C_4 : j \leq 9$
	$C_5 : C - 0.001 < 0$
<u>Equations</u>	
$E_1 : x_{1,j} = 1$	
$E_2 : x_{1,j} = 0$	
$E_3 : C = 0$	
$E_4 : E = (x_{i-1,j} + x_{i+1,j} + x_{i,j-1} + x_{i,j+1})/4$	
$C = C +  E - x_{i,j} $	
$x_{1,j} = E$	
$E_5 : STOP$	

$S_1$  sets up the mesh,  $S_2 C_5$  represents the main calculation, and  $E_5$  stops the problem after the iteration is complete. By defining the additional range statements:

$$I_3 : i = 1(9)10 \quad \text{and} \quad J_3 : j = 1(9)10$$

we could set up the mesh using no condition statements whatever:

$S_1 : I_3(J_1 E_1) + I_2(J_3 E_1) + I_2(J_2 E_2)$  which would probably improve the set-up from a time standpoint.

Note that the use of C and E as program variables is not restricted, even though the same symbols are used for control.

### 3. Algae Programming for the 704.

#### 3.1 Basic Assumptions and Restrictions.

Since Algae I currently uses the Fortran compiler as an intermediate step in its assembly, it is assumed that the programmer has a sound knowledge of basic Fortran concepts and restrictions. In addition, the following rules apply:

1. All subscripts on the logical symbols (E,S,T,G, etc.) must be numerical and greater than zero. Subscripts on non-control symbols are restricted only by the Fortran regulations.
2. Formula numbers can be assigned by the Algae programmer only to FORMAT statements, and must lie within the range  $1 \leq FN \leq 99$ . All other formula numbering will be done by the Algae code under one of two options selected by the programmer.
3. Because of restriction 2, FREQUENCY statements cannot be used in Algae I.
4. Comments can appear only at the beginning of the deck or within the defined E's.
5. Identification punched in columns 73-80 will be reproduced on Algae I listings only if the cards are read onto tape by the off-line equipment.

#### 3.2 Coding Conventions and Specific Restrictions.

The Fortran coding form is used to prepare the problem for keypunching. Specific instructions for defining each character are given below with the restrictions governing the use of this character in Algae I. Unless otherwise defined, subscripts must lie in the range  $1 \leq i \leq 999$ . There may be no more than 500 E's, each having 70 cards or less.

E<sub>i</sub> Each E<sub>i</sub> is a list, in Fortran language, of the "Things to be done" by one flow diagram "box" or closed set of equations. None of the statements listed within an E are tested by Algae I for accuracy. Rather, Algae I substitutes the entire set of equations into the sequence of Fortran statements whenever the specific E is requested by a control statement.

E's may contain comments at any point (they may consist entirely of comments), input-output statements, and/or the two control statements SENSE LIGHT i and PAUSE.\* These are

\*The control statement STOP should not be used, since it results in a logic error in FORTRAN.

the only control statements listed in chapter 4 of the Fortran manual that should ever be used by the Algae programmer. E's cannot contain FORMAT, DIMENSION, or EQUIVALENCE statements, and must be less than 70 cards in length. Multiple reference may be made to defined E's, as described in section 2.2.

Page 14 illustrates the method of defining E's. Note that the definition of E<sub>32</sub> is terminated by the appearance of another identification symbol, E<sub>16</sub>. E<sub>16</sub> is in turn terminated by the symbol C<sub>3</sub>. Fortran equations within an E may be continued from one card to the next as in normal Fortran notation.

C<sub>1</sub>

There are two types of conditions in the Algae I program: Type I, illustrated by C<sub>3</sub> and C<sub>5</sub> on page 14 defines the test of a numerical condition. The term or expression within parentheses is compared to zero in a manner designated by the symbol appearing at the right of the parenthesis. If we denote the contents of the parenthesis by X,

(X) P asks, "Is X positive?"	(X > 0?)
(X) N asks, "Is X negative?"	(X < 0?)
(X) = asks, "Is X equal to zero?"	(X = 0?).

The symbols may be negated to obtain the remaining inequalities:

(X) P <sup>*</sup> , "Is X ≤ 0?"
(X) N <sup>*</sup> , "Is X ≥ 0?"
(X) = <sup>*</sup> , "Is X ≠ 0?"

In Algae I, the numerical condition (including parentheses) must occupy at most 48 card columns.

Type II conditions, illustrated by C<sub>1</sub>, C<sub>2</sub>, and C<sub>4</sub> on page 14, define the test of a trigger, sense switch, or sense light. It is distinguished from a type I condition by the absence of a left parenthesis as the initial non-blank character. A comparator is also used to indicate whether the condition is satisfied when the trigger is on or off.

Typical type II conditions are:

SENSE LIGHT 3 ON	S L 3 ON
SENSE SWITCH 6 OFF	SW 6 OFF
ACCUMULATOR OVERFLOW ON	A ON
DIVIDE CHECK ON	D ON
QUOTIENT OVERFLOW OFF	Q OFF

For the benefit of programmers who are in a hurry or dislike spelling, Algae I will accept the abbreviations at the right in the above listing. An S, A, D, or Q must appear in the first non-blank position to avoid an error, and the ON/OFF condition must be specified. Do not attempt to abbreviate the SENSE LIGHT i instruction when it is used within an E to turn a specific light on.

A  $C_i$  of the two above types must have a numerical subscript within the range  $1 \leq i \leq 149$ .

C-Statements, illustrated by  $C_{153}$  on page 14 and defined in section 1.5, enable the programmer to use repeatedly a complex condition with a minimum of writing. A C-Statement must have as its identification a numerically subscripted  $C_i$  with  $150 \leq i \leq 200$ . In Algae I, a C-Statement cannot contain another C-Statement, and must be complete on one card (no continuations).

A C Group has been defined as any closed grouping of conditions, including condition statements. Thus if a statement contains the expression  $\dots + E_3 C_4 + C_1 (C_2 + C_3) * C_{155} E_2 + \dots$ ,  $C_4$  and  $C_1 (C_2 + C_3) * C_{155}$  are C groups. The maximum number of C's which may appear in a C group is 70. In the second group above, this would include all C's within the statement  $C_{155}$  but not  $C_{155}$  itself.

$I_i - N_i$ : Range statements (loop conditions) are identified by the numerically subscripted characters I, J, K, L, M, N. Examples of range statements are  $I_4$  and  $M_2$  on page 14. All range statements have the form  $a=b(c)d$  where:

a is the subscript whose range is being defined,

b is the first value to be taken by a,

c is the increment of change, and

d is the final bound on a. The final value of a is  $\leq d$  if  $c > 0$ , or  $\geq d$  if  $c < 0$ .

The following rules apply:

1. a may consist of any four Hollerith characters, the first of which is I, J, K, L, M, or N.
2. b and d may be any non-zero symbols recognized by Fortran as fixed point constants or variables. They can never assume negative values. Note that Fortran prevents the use of subscripted fixed point variables for b and d in DO's.
3. c may be any non-zero symbol (positive or negative) recognized by Fortran as a fixed point constant or variable. If d is greater than b, c must be greater than zero. If b is greater than d, Algae I will set up a reverse DO loop using a dummy variable. In this case, c must be preceded by a minus sign to indicate the reversal of the loop condition. The dummy variable is formed by repeating the first letter of a three times and adding the remainder of the term (see below). Care should be taken to avoid using this variable in the problem.

Some acceptable forms would be:

IVAL = IMAX(-IDELTA)IMIN	Dummy Variable: IIIIVAL
J32 = 45 (-3) JLOW	Dummy Variable: JJJ32
M3K2 = 1 (INCR) 10	
L = L1 (3) 15.	

It is not necessary for the symbol a to be the same as the identification symbol. For example, in  $I_4$  on page 14 we define the range of the subscript JT. The subscript on the identification must lie in the range  $1 \leq i \leq 999$ . Only 150 range statements can be defined in a problem. Range statements, once defined, may be used repeatedly.

All 4 fields (a, b, c, and d) must be filled. Algae does not make the assumption that a missing term is equal to 1.

S<sub>1</sub> and T<sub>1</sub> These control statements are written for the 704 in the same manner as those defined in examples 1, 1a, and 2 (see page 14 for further illustrations). Note that all subscripts are numerical and non-zero. Subscripts may be written as

smaller characters, but caution should be taken to keep them above the line to avoid keypunching errors. Care should also be taken to make commas distinctively different from the subscript 1.

S's and T's may be continued over several cards without the use of a continuation symbol in column 6, however a statement must have less than 1000 characters (not counting subscripts) and must have any subscripts completed on one line. One problem may contain a maximum of 450 S's and T's combined.

Subscripts on  $S_i$  must lie in the range  $1 \leq i \leq 999$  but subscripts on  $T_j$  are restricted to the range  $1 \leq j \leq 100$ . To provide the multiple exit to a T sequence, the dummy variable  $LL_{xxx}$  is formed, where  $xxx$  is the subscript on the defined T. Thus,  $T_{56}$  would produce a dummy variable  $LL_{56}$  which should not be used elsewhere in the code.

$G_i$  and  $H_i$  These characters need never be defined and should appear only in S or T statements. In both cases the subscript  $i$  should lie between  $1 \leq i \leq 200$ . Violation of this restriction will not produce an error stop, but might lead to a swifter depletion of the allowed quantity of formula numbers.

### 3.3 Deck Ordering and Operation.

#### 1. Deck ordering.

A problem designed for Algae I assembly may be considered to be divided into two sections. Section I must contain all

1. Control cards (see part 2 below)
2. Initial comments (all other comments must appear within E's)
3. Format statements
4. Dimension statements
5. Equivalence statements.

These cards may be arranged in any order in section I. All cards in section I must be loaded before the first card of section II. Section II consists of all the defined E's, S's, T's, C's and range statements, in any order.

#### 2. Sense switches.

Sense switches 5 and 6 are used by the 8K ALGAE I. They have the following effect:

##### Sense switch 5:

If this switch is down, and if errors have been detected by the compiler, the FORTRAN code, if any, which has been produced will be listed on tape 9.

##### Sense switch 6:

If this switch is up, each formula in an E will be assigned a formula number. If the switch is down, formula numbers will be assigned only where necessary to define control of the problem.

#### 3. Preparation of Algae System Tape.

The Algae System Tape is tape # 6. Because of the large number of cards involved, Algae should not be read by the card reader unless the installation does not possess enough tape units to permit the entire system to function from tapes. To write the system tape, place a blank tape on logical tape 6 and load the ALGAE 8K deck in the card reader. Press "LOAD CARDS." Note that it is not necessary to clear the machine or reset the console.

WARNING: DO NOT ATTEMPT TO WRITE TWO TAPES SIMULTANEOUSLY BY TURNING THEM BOTH TO 6. THE WRITING PROGRAM READS THE TAPE IT HAS WRITTEN, AND COMPARES IT WITH THE CONTENTS OF MEMORY. The 704 cannot read 2 tapes at once, and the program will continue to try to write a good tape indefinitely.

#### 4. Operation

Input may be from cards or from a tape prepared on peripheral card-to-tape equipment. Each problem must be loaded separately, whether it is loaded by cards or by tape, as the input tape is destroyed in the process of compiling.

Machine set-up is as follows:

- a) Tapes by logical tape number.
  1. FORTRAN
  2. blank
  3. blank or BCD input (see above)
  4. blank
  5. blank
  6. ALGAE
  9. tape-to-printer BCD tape
- b) Printer. Share # 2 board.
- c) Punch. Blank cards.
- d) Card reader.
  1. ALGAE RREC 1 card.
  2. Problem decimal deck, if input is not on tape 3. Otherwise, no cards.
  3. No cards.

To initiate compiling, press the "LOAD CARDS" button. When the deck has been completely read in, and the heading printed on the printer, place ALGAE RREC 2 in the card reader and run it in to ready. If it is desired, the next deck to be compiled may follow RREC 2 in the card reader.

There are three program stops.

1. If a card being read in has a non-Hollerith character, the program will stop with  $120_8$  in the SR. The card which has just been read should be corrected, it and all unread cards replace in the card reader, and the "START" button pushed.
2. If compiling halts because of an error, a stop will occur with 5 in the address of the SR. Press "START" to load a self-loading card.
3. If there has been a machine error, there will be a stop with  $1747_8$  in the SR. Do not proceed until the trouble has been determined.



Note: The RREC 2 card brings into memory a code which transfers the FORTRAN output from tape 2 to tape 9. If the local FORTRAN has such a routine, do not use RREC 2.

Note: As soon as RREC 2 has been read, a new tape 3 with new BCD input may be hung.

C ← comment	For continuation	Statement number	Code	Text	Remarks
1	5	72	73	80	Identification
		1		7	80
				EXAMPLES OF ALGAE I DEFINITIONS:	
		C E 32		COMMENTS MAY APPEAR ANYWHERE WITHIN AN E	REMARKS
				X = X**2 + 3.2 * Y + SQRTF(VM)	MAY BE
				T(I) = U(NP(S))**2 + V(NMIN) + W(N3) - SQRTF(Z**3) / 4.0*N1	PUNCHED
				x + N2	HERE -
		E 16		PAUSE	PRINTED
				BE SURE TO IDENTIFY EACH COMMENT BY A C IN	DIFF - LINE
				C COLUMN 1.	BY ALGAE
				(X**2 - 3.0*Y(I) + 7.29) P*	
				(Y-2) =	
		C 1		DIVIDE CHECK 0 N	
		C 2		SENSE LIGHT 3 0 FFF	
		C 4		S SW 4 0 0 N	
		C 153		C1 + (C3 C4* + C5) C6*	
		I 4		JT = 3(15) 4	
		M 2		M3 = MAX(-3) MIN	
		S 4		C1(E3 + E4 C2 + I4 (M2 (E6 + I3 E4)) + T3 + C1 (S5) S6) + J3 (E4 + I2 (C1 C6 E7 + S3))) + E16	
		T 71		K2 (E1 + (E2, E3) C4) + C1 G2	
				(The above statements are not intended to form a consistent set.)	

### 3.4 Sample Problems

On the following pages will be found two sample problems coded in Algae. The first problem is correctly coded. Most questions arising about the use of forms, control cards, etc., may be answered by careful perusal of this listing.

In the second code, "Laplace", on page 25 ff., intentional errors have been made so that the reader may become familiar with the diagnostic procedures used in Algae I. For a more complete discussion of diagnostics, see section 4. of this writeup.

```

ALG C ALGAE DEMONSTRATION CODE
C ARITHMETIC-GEOMETRIC MEAN COMPUTER
DIMENSION XP(25), YP(25), X(25), Y(25), FXY(25, 25)
EQUIVALENCE (X, XP), (Y, YP)
FORMAT (2I3)
2 FORMAT (35 H1 ALGAE DEMONSTRATION CODE / 8 H0 / 45 H
ARITHMETIC-GEOMETRIC MEAN COMPUTER / 8 H0 I
B= 13, 4 H, J= 13)
3 FORMAT (36 H0 X(I) Y(J) F(I,J) )
4 FORMAT (3 E 13.4)
5 FORMAT (36 H0 F(I,J) FROM PREVIOUS ITERATION / (8
AE 13.4))
6 FORMAT (35 H0 SUMX SUMY AGM / 3E13.4)
CE1 COMPUTE 2 PI. SET UP DXP, DYP.
CAP1 = ICAP
CAPJ = JCAP
TWOPI = 2.0 * 3.14159
DX = TWOPI / CAP1
DY = TWOPI / CAPJ
T = I-1
E23 COMPUTE THE VALUES OF X PRIME.
C XP(I) = T * DX
E24 COMPUTE THE VALUES OF Y PRIME.
C YP(J) = T * DY
E3 AM = AM1
GM = GM1
E4 X(I) = SIN(XP(I))
C COMPUTATION OF X.
CE5 FIRST COMPUTATION OF Y.
Y(J) = COS(EXPF(YP(J)))
CE6 SECOND COMPUTATION OF Y.
Y(J) = COS(EXPF(-YP(J)))
E7 AM = X(I)
C PREPARE AGM COMPUTATION.
GM = Y(J)
CE8 AGM COMPUTATION.
C ARITHMETIC MEAN.
AM1 = (AM + GM) / 2.0
C GEOMETRIC MEAN.
GM1 = SQRT(AM * GM)
E2 AGM = AM1
E9 FXY(I,J) = AGM
F10 FXY(I,J) = 0.0
E11 FXY(I,J) = -1.0
E12 FXY(I,J) = 1.0
CF13 INITIALIZE SUMS.
SUMX = 0.
SUMY = 0.
E14 SUMX = SUMX + ABSF(XP(I))
E22 PRINT 6, SUMX, SUMY, AGM
E15 SUMY = SUMY + ABSF(YP(J))
CE16 SET TO COMPUTE AGM OF SUMS.
AM = SUMX
GM = SUMY
SENSE LIGHT 1
E18 READ 1, ICAP, JCAP
PRINT 2, ICAP, JCAP
E19 PRINT 4, X(I), Y(J), FXY(I,J)
E20 PRINT 3
E21 PRINT 5, ((FXY(I,J)), I = 1, ICAP), J = 1, JCAP)
E25 SENSE LIGHT 2
CE26 IGNORE TEST.
E001.
E001.01
E001.02
E001.03
E001.04
E001.05
E023.01
E023.
E023.02
E024.01
E024.
E024.02
E003.01
E003.02
E004.01
E004.
E005.
E005.01
E006.
E006.01
E007.01
E007.
E007.02
E008.
E008.01
E008.
E008.02
E002.01
E009.01
E010.01
E011.01
E012.01
E013.
E013.01
E013.02
E014.01
E022.01
E015.01
E016.
E016.01
E016.02
E016.03
E018.01
E018.02
E019.01
E020.01
E021.01
E025.01
E026.

```

N1 J = ICAP (-1) 1  
 M2 J = 1 (1) JCAP  
 C1 (X(I)) =  
 C2 (Y(J)) =  
 C3 (X(I)) N  
 C4 (Y(J)) N\*  
 C5 SWITCH 6 OFF  
 C7 ACCUM OVERFLOW ON  
 C151 (C1 C4 C2\* + C2 C3\* C1\*)  
 C8 S LITE 1 OFF  
 C9 (ABS F (AM1 - GM1) / GM1 - .0001) N  
 C10 SENSE LT 2 ON  
 T1 (E3, E8) C9 + E2  
 T2 E1 + N1 E23 + M2 E24  
 S1 C3\* C1\* C4 C2\* (E7 + T1 + C7 (E10, E9), C151 (E12, C3 C4\*  
 (E11, E10)))  
 S2 T2 + N1 E4 + C8 (M2 E5, M2 E6) + C5\* E20  
 +N1 (M2 (S1 + C5\* E19)) + E21  
 S3 C7 E26 + H1 + E18 + H2 + S2 + C10\* (S4, G1)  
 S4 E25 + T2 + E13 + N1 E14 + M2 E15 + E16 + T1 + F22 + G2

ALGAE OBJECT CODE FROM FORTRAN TAPE 2

```

C ALGAE DEMONSTRATION CODE
C ARITHMETIC-GEOMETRIC MEAN COMPUTER
DIMENSION XP(25), YP(25), X(25), Y(25), FXY(25, 25)
EQUIVALENCE (X, XP), (Y, YP)
1 FORMAT (2I3)
2 FORMAT (35 H1 ALGAE DEMONSTRATION CODE / 8 H0 I 45 H I
   ARITHMETIC-GEOMETRIC MEAN COMPUTER /
   B= 13, 4 H, J= 13)
3 FORMAT (36 H0 X(I) Y(J) F(I,J) )
4 FORMAT (3 E 13.4)
5 FORMAT (36 H0 F(I,J) FROM PREVIOUS ITERATION / ( 8
   AE 13.4))
6 FORMAT (35 H0 SUMX SUMY AGM / 3E13.4)
IF ACCUMULATOR OVERFLOW 0144, 0145 C007
0144 CONTINUE E026.
C IGNORE TEST.
0145 CONTINUE
0112 CONTINUE
READ 1, ICAP, JCAP
PRINT 2, ICAP, JCAP
0113 CONTINUE
LL002=0002
GO TO 0111
0132 CONTINUE
D00133 III = 1, ICAP, 1
I= 1+ ICAP-III
X(I) = SIN(XP(I))
COMPUTATION OF X.
0133 CONTINUE
0134 IF (SENSE LIGHT 1)
D00129 J= 1, JCAP, 1
FIRST COMPUTATION OF Y.
Y(J) = COS(XP(J))
0129 CONTINUE
GO TO 0136
0135 CONTINUE
D00130 J= 1, JCAP, 1
SECOND COMPUTATION OF Y.
Y(J) = COS(XP(J))
0130 CONTINUE
0136 CONTINUE
0137 IF (SENSE SWITCH 6)
PRINT 3
0138 CONTINUE
D00139 III = 1, ICAP, 1
I= 1+ ICAP-III
J= 1, JCAP, 1
D00131 IF(X(I))
0102 IF(X(I))
0101 IF(Y(J))
0100 IF(Y(J))
0124 CONTINUE
AM = X(I)
PREPARE AGM COMPUTATION.
GM = Y(J)
LL001=0002
GO TO 0110
0117 CONTINUE
IF ACCUMULATOR OVERFLOW
0118 CONTINUE

```

```

      FXY(I,J) = 0.0
      GO TO 0120
0119 CONTINUE
      FXY(I,J) = AGM
0120 CONTINUE
      GO TO 0126
0125 CONTINUE
      IF(X(I))
0107 IF(Y(J))
0106 IF(Y(J))
0105 IF(Y(J))
0104 IF(X(I))
0103 IF(X(I))
0121 CONTINUE
      FXY(I,J) = 1.0
      GO TO 0123
0122 CONTINUE
      IF(X(I))
0108 IF(Y(J))
0114 CONTINUE
      FXY(I,J) = -1.0
      GO TO 0116
0115 CONTINUE
      FXY(I,J) = 0.0
0116 CONTINUE
0123 CONTINUE
0126 CONTINUE
      IF (SENSE SWITCH 6)
0127 CONTINUE
      PRINT 4, X(I), Y(J), FXY(I,J)
0128 CONTINUE
0131 CONTINUE
0139 CONTINUE
      PRINT 5, ((FXY(I,J), I = 1, ICAP), J = 1, JCAP)
      IF (SENSE LIGHT 2)
0146 CONTINUE
      SENSE LIGHT 2
      LL002=0001
      GO TO 0111
0140 CONTINUE
      C INITIALIZE SUMS.
      SUMX = 0.
      SUMY = 0.
      DO0141 III = 1, ICAP, 1
      I = 1+ ICAP-III
      SUMX = SUMX + ABSF (XP(I))
0141 CONTINUE
      DO0142 J= 1, JCAP, 1
      SUMY = SUMY + ABSF (YP(J))
0142 CONTINUE
      C SET TO COMPUTE AGM OF SUMS.
      AM = SUMX
      GM = SUMY
      SENSE LIGHT 1
      LL001=0001
      GO TO 0110
0143 CONTINUE
      PRINT 6, SUMX, SUMY, AGM
      GO TO 0113
0147 CONTINUE
      GO TO 0112
0111 CONTINUE
      C COMPUTE 2 PI. SET UP DXP, DYP.

```

```

E010.01
E009.01
0105, 0107, 0105 C001
0105, 0106, 0106 C004
0121, 0105, 0121 C002
0122, 0104, 0122 C002
0122, 0103, 0103 C003
0121, 0122, 0121 C001
E012.01
0108, 0115, 0115 C003
0114, 0115, 0115 C004
E011.01
E010.01
0127, 0128 C005
E019.01
E021.01
0147, 0146 C010
E025.01
T002
E013.
E013.01
E013.02
N001
N001
E014.01
M002
E015.01
E016.
E016.01
E016.02
E016.03
T001
E022.01
T002
E001.

```

```

      CAPI = ICAP
      CAPJ = JCAP
      TWOPI = 2.0 * 3.14159
      DX = TWOPI / CAPI
      DY = TWOPI / CAPJ
      DOO149      III = 1, ICAP, 1
                  I = 1+ ICAP-III
      T = I-1
      C COMPUTE THE VALUES OF X PRIME.
      XP(I) = T * DX
0149 CONTINUE
      DOO150      J = 1, JCAP, 1
      T = J-1
      C COMPUTE THE VALUES OF Y PRIME.
      YP(J) = T * DY
0150 CONTINUE
      GO TO(0140,0132), LL002
0110 CONTINUE
      GO TO 0151
0152 CONTINUE
      AM = AM1
      GM = GM1
0151 CONTINUE
      C AGM COMPUTATION.
      C ARITHMETIC MEAN.
      AM1 = (AM + GM) / 2.0
      C GEOMETRIC MEAN.
      GM1 = SQRTF (AM * GM)
      IF(ABSF (AM1 - GM1) / GM1 - .0001)
0153 CONTINUE
      AGM = AM1
      GO TO(0143,0117), LL001

```

```

E001.01
E001.02
E001.03
E001.04
E001.05
N001
N001
E023.01
E023.
E023.02
M002
E024.01
E024.
E024.02
T001
E003.01
E003.02
E008.
E008.
E008.01
E008.
E008.02
0153, 0152, 0152 C009
E002.01

```



STORAGE FOR VARIABLES APPEARING IN DIMENSION OR EQUIVALENCE SENTENCES

DEC OCT      DEC OCT      DEC OCT      DEC OCT      DEC OCT  
 FXY 32717 77715    XP 32767 77777    X 32767 77777    YP 32742 77746    Y 32742 77746

STORAGE FOR VARIABLES WHICH DO NOT APPEAR IN DIMENSION OR EQUIVALENCE SENTENCES

DEC OCT      DEC OCT      DEC OCT      DEC OCT      DEC OCT  
 I 32092 76534      I11 32091 76533      ICAP 32090 76532      GM 32089 76531      GM1 32088 76530  
 DY 32087 76527      DX 32086 76526      CAPJ 32085 76525      CAPI 32084 76524      AM 32083 76523  
 AM1 32082 76522      AGM 32081 76521      JCAP 32080 76520      J 32079 76517      LL001 32078 76516  
 LL002 32077 76515      SUMX 32076 76514      SUMY 32075 76513      T 32074 76512      TWOP1 32073 76511

EXTERNAL FORMULA NUMBERS WITH CORRESPONDING INTERNAL FORMULA NUMBERS AND OCTAL LOCATIONS

EFN	IFN	LOC	EFN	IFN	LOC	EFN	IFN	LOC	EFN	IFN	LOC
1	3	00000	2	4	00000	3	5	00000	4	6	00000
6	8	00000	144	10	00002	145	11	00002	112	12	00003
132	22	00044	133	26	00064	134	28	00072	129	31	00103
130	36	00117	136	37	00121	137	39	00123	138	42	00134
101	48	00166	100	49	00172	124	50	00174	117	55	00207
119	60	00216	120	62	00220	125	64	00223	107	66	00227
105	68	00241	104	69	00247	103	70	00253	121	71	00255
108	76	00265	114	77	00270	115	80	00273	116	82	00275
126	84	00277	127	86	00301	128	90	00317	131	91	00317
146	102	00374	140	106	00402	141	112	00425	142	115	00437
147	126	00474	111	128	00476	149	138	00545	150	142	00566
152	146	00577	151	149	00603	153	153	00626			

SUBROUTINES OBTAINED FROM LIBRARY

DEC	OCT	DEC	OCT	DEC	OCT	DEC	OCT	DEC	OCT
(DBC)	491 00753	(CSH)	1008 01760	(BDC)	1089 02101	(FIL)	1119 02137	(LEV)	1132 02154
(RTN)	1159 02207	(SPH)	1639 03147	SORT	1731 03303	COS	1753 03331	SIN	1756 03334
EXP	1810 03422								

9A	TOV 10A	TXL 24A*4	STD 91A2
10A	TRA 11A	MSE 97	LXD 1*2
11A	BSS	TRA 28A	SXD C1100*2
D1403	LXD C1G0*4	TRA 33A	LXD C1100*1
12A	BSS	BSS	LXD C1G1*2
13A	CAL *	LXD 2)+1*4	CLA X+1*4
	XIT (LEV)	CLA JCAP	TZE 47A
	ETM	STD 31A2	TPL 47A
	CAL (DBC)	CLA YP+1*4	TRA 64A
	SLW 1	SXD 6)+4*4	CLA X+1*4
	CAL (CSH)	TSX EXP*4	TZE E120R
13D1	NTR 8)1*0*81	TSX COS*4	CLA Y+1*2
14A	ETM	LXD 6)+4*4	TZE 49A
	NTR ICAP	STO Y+1*4	TPL 49A
	NTR JCAP	BSS	TRA E110R
	LTM	TXI *+1*4*1	CLA Y+1*2
15A	CAL *	TXL 30A*4	TZE 64A
16A	XIT (RTN)	TRA 37A	BSS
	CAL *	BSS	CLA X+1*4
	XIT (LEV)	LXD 2)+1*4	STO AM
	ETM	CLA JCAP	CLA Y+1*2
	CAL (RDC)	STD 36A2	STO GM
	SLW 1	CLS YP+1*4	CLA 2)
	CAL (SPH)	SXD 6)+4*4	STO L1001
16D1	NTR 8)2	TSX EXP*4	LXD L1001*4
17A	ETM	TSX COS*4	SXD C1G3*4
	NTR ICAP	LXD 6)+4*4	PXD 0*2
	NTR JCAP	STO Y+1*4	STO J
	LTM	BSS	TRA 144A
18A	CAL *	TXI *+1*4*1	BSS
D1504	XIT (FIL)	TXL 35A*4	TOV 57A
D1404	LXD C1100*1	BSS	TRA EJP
19A	LXD C1G3*4	PSE 118	BSS
20A	RSS	TRA 42A	CLA 3)
	CAL 2)	BSS	STO FXY+1*1
	STO L1002	CAL *	TRA 62A
	LXD L1002*2	XIT (LEV)	SXD C1G1*2
	SXD C1G2*2	ETM	SXD C1100*1
21A	TRA 128A	CAL (RDC)	BSS
22A	BSS	SLW 1	CLA AGM
23A	LXD 2)+1*4	CAL (SPH)	STO FXY+1*1
	CLA ICAP	NTR 8)3	BSS
	STD 26A2	CAL *	TRA D1418
	PXD 0*4	XIT (FIL)	SXD C1G1*2
	STO III	BSS	SXD C1100*1
24A	CLA 2)+1	LXD 2)+1*4	SYN E110R
	ADD ICAP	SXD C)2*1*4	BSS
	SUB III	CLA ICAP	CLA X+1*4
	STO I	STD 92A2	TZE 66A
	LXD 1*2	PXD 0*4	TPL 68A
	SXD C1G0*2	STO III	TRA 68A
25A	CLA XP+1*2	CLA 2)+1	CLA Y+1*2
	SXD 6)+4*4	ADD ICAP	TZE E1T
	TSX SIN*4	SUB III	TPL E1T
	LXD 6)+4*4	STO I	TRA 68A
	STO X+1*2	LXD 1*4	SXD C1G1*2
26A	BSS	SXD C1G0*4	SXD C1100*1
26A1	TXI *+1*4*1	LXD 2)+1*1	CLA Y+1*2
	SXD III*4	SXD C1G1*1	TZE 68A
		CLA JCAP	TPL 71A

68A TRA 71A  
 68A1 CLA Y+1\*2  
 TZE EJV  
 TPL 74A  
 TRA 74A  
 EIV SXD CIG1\*2  
 SXD C1100\*1  
 69A CLA X+1\*4  
 69A1 TPL 70A  
 TPL 70A  
 TRA 74A  
 70A CLA X+1\*4  
 70A1 TZE 74A  
 71A BSS  
 72A STO FXY+1\*1  
 73A TRA 83A  
 74A BSS  
 75A CLA X+1\*4  
 75A1 TPL 80A  
 TPL 80A  
 E113 SXD CIG1\*2  
 SXD C1100\*1  
 76A CLA Y+1\*2  
 76A1 TZE 80A  
 TPL 80A  
 77A BSS  
 78A CLS 3)+1  
 79A STO FXY+1\*1  
 TRA 82A  
 80A BSS  
 81A CLA 3)  
 STO FXY+1\*1  
 82A BSS  
 83A BSS  
 TRA 84A  
 D1118 LXD CIG0\*4  
 84A BSS  
 85A PSE 118  
 TRA 90A  
 86A BSS  
 87A CAL \*  
 XIT (LEV)  
 ETM  
 CAL (BDC)  
 SLW 1  
 CAL (SPH)  
 NTR 8)4  
 87D1 ETM  
 88A NTR X+1\*4  
 NTR Y+1\*2  
 NTR FXY+1\*1  
 LTM  
 CAL \*  
 XIT (FIL)  
 BSS  
 90A BSS  
 91A BSS  
 91A1 TXI \*\*1,2\*1  
 TXI \*\*1,1\*25  
 91A2 TXL 46A\*2

92A BSS C1201\*2  
 LXI \*\*1,2\*1  
 92A1 SXD C1201\*2  
 SXD C1201\*2  
 92A2 TXL 111\*2  
 93A CAL \*  
 TXL 44A\*2  
 CAL \*  
 XIT (LEV)  
 ETM  
 CAL (BDC)  
 SLW 1  
 CAL (SPH)  
 93D1 NTR 8)5  
 94A LXD 2)+1\*2  
 CLA ICAP  
 STD 97A2  
 LDQ JCAP  
 MPY 2)+2  
 ALS 17  
 STD 98A4  
 CLA ICAP  
 SUR 6)+3  
 STD 98A3  
 ADD 2)+2  
 STD 98A1  
 CLA ICAP  
 STD 97A3  
 95A BSS  
 96A BSS  
 97A ETM  
 NTR FXY+1\*2  
 LTM  
 TXI \*\*1,2\*1  
 97A1 TXL 97A\*2  
 97A2 TIX 97A3+1\*2  
 97A3 TXI 98A1+1\*2  
 98A1 SXD 97A2\*2  
 98A3 TIX 98A3+1\*2  
 98A4 TXL 95A\*2  
 99A LTM  
 100A CAL \*  
 XIT (FIL)  
 MSE 98  
 101A TRA 102A  
 TRA 126A  
 BSS  
 102A CLA 2)+1  
 103A STO L1002  
 104A LXD L1002\*2  
 SXD CIG2\*2  
 TRA D111P  
 BSS  
 105A CLA 3)  
 106A STO SUMX  
 107A CLA 3)  
 108A CLA 3)  
 109A STO SUMY  
 LXD 2)+1\*4  
 CLA ICAP  
 STD 112A2

PXD 0\*4  
 STO III  
 110A CLA 2)+1  
 ADD ICAP  
 SUB III  
 STO 1  
 LXI 1\*2  
 111A SXD CIG0\*2  
 CLA XP+1\*2  
 SSP  
 FAD SUMX  
 STO SUMX  
 112A BSS  
 112A1 TXI \*\*1,4\*1  
 SXD III\*4  
 112A2 TXL 11CA\*4  
 113A LXI 2)+1\*4  
 CLA JCAP  
 STD 115A2  
 114A CLA YP+1\*4  
 SSP  
 FAD SUMY  
 STO SUMY  
 115A BSS  
 115A1 TXI \*\*1,4\*1  
 115A2 TXL 114A\*4  
 116A CLA SUMX  
 STO AM  
 117A CLA SUMY  
 STO GM  
 118A PSE 97  
 119A CLA 2)+1  
 STO LL001  
 LXI LL001\*4  
 SXD CIG3\*4  
 120A TRA D131U  
 E11N SXD CIG1\*2  
 SXD C1100\*1  
 BSS  
 121A CAL \*  
 122A XIT (LEV)  
 ETM  
 CAL (BDC)  
 SLW 1  
 CAL (SPH)  
 NTR 8)6  
 122D1 ETM  
 123A NTR SUMX  
 NTR SUMY  
 NTR AGM  
 LTM  
 CAL \*  
 XIT (FIL)  
 TRA 19A  
 BSS  
 124A BSS  
 125A TRA 12A  
 126A LXI C1100\*1  
 D111P BSS  
 127A BSS  
 128A CLA ICAP  
 129A LRS 18

130A	ORA 6) FAD 6) STO CAPI CLA JCAP LRS 18	D)31U D)21U 144A 145A 146A 147A 148A 149A 150A	TRA 22A TRA 106A LXD C)10,1 LXD C)G1,2 BSS TRA 149A BSS CLA AM1 STO AM CLA GM1 STO GM BSS CLA AM FAD GM FDP 3)+2 STO AM1 LDQ AM FMP GM	BCD 10DE BCD 1 / BCD 1 45 H BCD 1 BCD 1 A BCD 1RITHME BCD 1TIC-GE BCD 1OMETRI BCD 1C MEAN BCD 1 COMPU BCD 1TER BCD 1 / BCD 1 8 HO BCD 1 I BCD 1= I3 BCD 1, 4 BCD 1H, J= BCD 1 I3) OCT -377777777777 BCD 1 ( BCD 136 HO BCD 1 X( BCD 1I) BCD 1 Y BCD 1(J) BCD 1(J) BCD 1(J) BCD 1) OCT -377777777777 BCD 1 ( BCD 13 E 13 BCD 1*4) OCT -377777777777 BCD 1 ( BCD 136 HO BCD 1 F( BCD 1(J) F BCD 1(J) BCD 1) OCT -377777777777 BCD 1 ( BCD 13 E 13 BCD 1*4) OCT -377777777777 BCD 1 ( BCD 136 HO BCD 1 F( BCD 1(J) F BCD 1R0M PR BCD 1EVI0US BCD 1 ITERA BCD 1TION BCD 1 / BCD 1 (8 BCD 1E 13*4 BCD 1I) OCT -377777777777 BCD 1 ( BCD 135 HO BCD 1 SUM BCD 1X BCD 1X BCD 1MY BCD 1 BCD 1AGM BCD 1 / 3 BCD 1E13*4) OCT -377777777777	8)3	
131A	STO CAPJ LDQ 3)+2 FMP 3)+3	151A	FMP GM			
132A	STO TWOPI CLA TWOPI		SXD 6)+4,4 TSX SORT,4 LXD 6)+4,4			
133A	FDP CAPJ CLA TWOPI	152A	CLA AM1 FSB GM1 SSP FDP GM1 STO GM1 CLA AM1 FSB GM1			
134A	LXD 2)+1,4 CLA ICAP STD 138A2 PXD 0,4		LXD 6)+4,4 SXD 6)+4,4 TSX SORT,4 LXD 6)+4,4	8)4		
135A	STO III CLA 2)+1 ADD ICAP SUB III STO I LXD I,1 SXD C)G0,1		CLA AM1 FSB GM1 SSP FDP GM1 STO GM1 CLA AM1 FSB GM1	8)5		
136A	SUB 2)+1 LRS 18 ORA 6) FAD 6) STO T LDQ T FMP DX STO XP+1,1	152A1 153A 154A 155A	TZE 146A TPL 146A ESS CLA AM1 STO AGH TRA 155A+3,4 TRA 55A TRA E)IN OCT +000002000000 OCT +000001000000 OCT +000031000000 OCT +000000000000 OCT +201400000000 OCT +202400000000 OCT +202622077174 OCT +163643334272 OCT +233000000000 OCT +000000077777 OCT +000000000000 OCT +000001000000 OCT +000000000000 OCT +000000000000 BCD 1 BCD 1213) OCT -377777777777 BCD 1 BCD 135 H1 BCD 1 BCD 1 ALG BCD 1AE DEM BCD 1ONSTRA BCD 1TION C	8)6		
137A	STO XP+1,1 BSS TXI *+1,4,1 SXD III,4 TXL 135A,4 LXD 2)+1,4 CLA JCAP STD 142A2 PXD 0,4 STO J CLA J SUB 2)+1	2) 3) 6)				
138A 138A1	TXI *+1,4,1 SXD III,4					
138A2 139A	TXL 135A,4 LXD 2)+1,4 CLA JCAP STD 142A2 PXD 0,4					
140A	STO J CLA J SUB 2)+1	8)1 8)2				
141A	LDQ T FMP DY STO YP+1,4					
142A 142A1	BSS TXI *+1,4,1					
142A2 143A	SXD J,4 TXL 140A,4 TRA 143A+3,2					

```

ALG C   ALGAE LAPLACE EQUATION ROUTINE -- TYPE 1 MESH
C       THIS CODE SOLVES THE LAPLACE EQUATION ON A 10 X 10 MESH
C       WITH BOUNDARIES SET TO 1.
1       DIMENSION X(10, 10)
2       FORMAT (30H1ALGAE LAPLACE EQUATION SOLVER/ 43H 10 X 10 MESH,
3       A WITH BOUNDARIES SET TO 1)
C       FORMAT (10 E 12,4)
C       FORMAT (36H0ITERATION ON LAPLACE EQUATION MESH / (10E11,3)
C
C       THIS COMPILING WILL DEMONSTRATE ERROR FORMATS.
C
S2      E3 + I2 (J2 E)
S1      I (J1 ((C1 + C2 + (C3 C4)* ) (E1, E2))
S3      S1 + (S2 C6 E7) C5 + E5 + E8
I1      I = 1 (1)
I2      I = 2 (1) 9
J1      J = 1 (1) 10, AND SO ON
J2      J = 2 (1) 9
C1      (I-1)
C2      (I-10) = EQUAL
C3      (J-2) N*
C4      (J-9) P*
C5      C = 0.001) N
C6      (SW 6 ON)
CE1     SET UP BOUNDARY VALUES OF MESH
X(I,J) = 1.0
CE2     SET INTERIOR TO ZERO
X(I,J) = 0.0
E3      C = 0.0
CE4     CALCULATE DIFFERENCES,
E       E = (X(I-1,J) + X(I+1,J) + X (I,J-1) + (XI,J+1)) / 4.0
C       SUM DIFFERENCES,
C       C = C + ABSF (E - X(I,J))
C       AND SET NEW VALUE OF X(I,J).
X(I,J) = E
E5      PRINT 1
E6      PRINT 2, (X)
E7      PAUSE
        PRINT 3, (X)
S1      IMPROPER I 0
S1      HAS TOO MANY (,S
C1      IMPROPER STATEMENT
C2      IMPROPER STATEMENT E
C4      IMPROPER STATEMENT C
C5      IMPROPER STATEMENT C
C6      IMPROPER STATEMENT W0
S1      IMPROPER
S2      UNDEFINED E0
S3      HAS IMPROPER TERM E7
        UNDEFINED E8
CODE 07241
CODE 07042
CODE 10222
CODE 10254
CODE 10535
CODE 10346
CODE 10222
CODE 04715
CODE 04624
CODE 04451
CODE 04624

```

```

C ALGAE LAPLACE EQUATION ROUTINE --- TYPE 1 MESH
C THIS CODE SOLVES THE LAPLACE EQUATION ON A 10 X 10 MESH
C WITH BOUNDARIES SET TO 1.
1 FORMAT (30H1ALGAE LAPLACE EQUATION SOLVER/ 43H 10 X 10 MESH,
2 A WITH BOUNDARIES SET TO 1)
2 FORMAT (10 E 12.4)
3 FORMAT (36H0ITERATION ON LAPLACE EQUATION MESH / (10E11.3)
C
C THIS COMPILING WILL DEMONSTRATE ERROR FORMATS.
C
0107 CONTINUE
0104 CONTINUE
0102 C = 0.0
      D00103
      D00101
0101 CONTINUE
0103 CONTINUE
      IF(SW 6 ON)
0105 CONTINUE
0106 PRINT 3, (X)
      CODING ERROR
0108 CONTINUE
0109 PRINT 1
0110 PRINT 2, (X)
      I= 2, 9, 1
      J= 2, 9, 1
      E003.01
      I002
      J002
      0105, 0105, 0105 C006
      E007.01
      0107, 0107 C005
      E005.01
      E005.02
      FIX AND RE-TRY
      0011 DIAGNOSTICS
      HALT
      CODE 07563

```

#### 4. Error Procedure and Description of Diagnostics

Algae I is designed to detect errors in source programs presented to it, and to give as part of its output as much information as is possible concerning those errors.

The compiling process may be divided into two phases. In Phase 1, the source program is read into the 704 DPM, statements are checked for duplication and improper subscripts, and diverse tables needed in Phase 2 are set up. In Phase 2 the FORTRAN source program is compiled. Detection of an error in Phase 1 prevents the execution of Phase 2. Detection of an error in Phase 2 prevents the execution of FORTRAN.

Errors may be divided into four major classes. These are:

- I. Machine errors.
- II. Phase 1 source program errors.
- III. Phase 2 source program errors which are relatively limited in their effect upon the code.
- IV. Phase 2 source program errors whose nature makes further compiling impossible.

Class I errors may cause a program stop at 1747<sub>8</sub>. An identification number, ALPHA, is printed on the line printer together with a brief description of the failure. Class II, III, and IV errors cause a statement to be printed on the off-line listing, containing an identification number, ALPHA, and a brief description of the source program error. Class II errors will cause the next problem to be loaded at the end of Phase I. Class III errors will cause the next problem to be loaded at the end of Phase 2. Class IV errors will cause the next problem to be loaded whenever they are detected. Before leaving a problem the total number of source program errors which have been detected is printed both on the line printer and off-line listing.

Information printed about errors is listed in the following format: H ID COMMENT CHAR TYPE ALPHA

The H position contains the word HALT if processing of the problem is interrupted for any reason. ID is the identification of the objectionable statement when available. COMMENT is a brief description of the error. CHAR is the term of the statement, if any, to which exception is taken. TYPE indicates whether the error

is a machine failure or source program mistake. ALPHA is the octal location at which the error was detected. By looking up this number in the following ALPHA table, a slightly expanded description may be obtained.

There are two symbols which may occur in the CHAR column on the listing which are the result of coding errors, but which do not appear in the source program at any point.

The first of these is  $W_0$ . This character is inserted at one stage in the code when a character has been determined to be improper in some respect. The insertion of  $W_0$  permits Algol I to continue compiling, and also permits the programmer to trace the effect of his errors on the resulting code.

The second is an E with subscript  $\geq 1000$ . This symbol is the identification of the entire coded contents of a parenthetical expression. It sometimes happens that a program will give rise to a vacuous parenthesis level, or other similar error. Then, when it is desired to incorporate this grouping into the rest of the code, the  $E_{1000}$  symbol is printed.

The format of the ALPHA table is as follows:

ALPHA	COMMENT	CLASS of error
	Nature of error	What to do about it

ALL ALPHA's are in octal.



### 5.0 Summary on Referencing

The following table gives, in condensed form, most of the restrictions on the Algae I source program.

Statement Type	Subscript Range	max. no. defined	no. of times each may be used	$\frac{\text{no. of cards}}{\text{no. of char.}}$	Remarks
$S_i$	$1 \leq i \leq 999$	500	1	$\frac{70}{1000}$	Control statement
$T_i$	$1 \leq i \leq 100$	100	100	$\frac{70}{1000}$	Subroutine
$C_i$	$1 \leq i \leq 149$	149	Indefinite	$\frac{1}{48 \text{ including blanks \& } ( )}$	Simple Condition
$C_i$	$150 \leq i \leq 200$	51	Indefinite	$\frac{1}{66}$	C-statement. May not contain another C-statement.
$I_i - N_i$	$1 \leq i \leq 999$	150	Unlimited	$\frac{1}{\text{See p. 10}}$	Range statement
$G_i$	$1 \leq i \leq 200$	$200^{\dagger}$	Unlimited	$\frac{0}{0}$	Transfer statement
$H_i$	$1 \leq i \leq 200$	$200^{\dagger}$	1	$\frac{0}{0}$	Point of transfer entry.

$^{\dagger}i$  may be greater than 200, but in large problems the available formula numbers may be exhausted.