

**GENIE**

# GENIE

Genie	.....
Genie Program Format	.....
Names	.....
Numbers	.....
Variables	.....
Declarations	.....
Functions	.....
Constants	.....
Remarks	.....
Command Sequence	.....
Arithmetic Expressions	.....
Arithmetic Commands	.....
Conditional Arithmetic Commands	.....
Transfer Control Commands	.....
Loop Control Commands	.....
Storage Control Commands	.....
Execute Control Commands	.....

## GENIE (continued)

Input-Output Commands . . . . .	
(Including Sense lights)	
Data Commands . . . . .	
Fast Registers . . . . .	
Assembly Language . . . . .	
Punctuation . . . . .	
Compilation Procedure . . . . .	
Running Genie Programs . . . . .	
Coding Examples . . . . .	
Genie Coding Conventions . . . . .	

The formula language for the Rice Computer is called the Genie language. Programs written in the Genie language are called Genie programs. Translation of Genie programs into machine language is accomplished by the Genie compiler.

The language and the compiler are both often referred to as just Genie. What is meant is usually clear from the context.

Genie programs may contain instructions written in the AP2 assembly language. Hence, the AP2 assembly language is a subset of the Genie language, and the AP2 assembly program is a subset of the Genie compiler.

## PROGRAM FORMAT

The unit of definition to the Genie compiler is the definition set, which has the form

```

    DEFINE
        declarations of external variables and
        parameters for the entire definition set
        constant codeword address specifications for external
        variables
        function definitions
    PROG1(PARAM1).=SEQ
        declarations of internal variables
        remarks
        constant specifications
        command sequence for the calculation
    END
    PROG2(PARAM2).=SEQ
        :
        :
    END
        :
        :
    PROGN
        :
        :
    END
    DEFINE
    LEAVE
    | cr      | 1st tab stop

```

1<sup>st</sup> program  
in definition  
set  
  
 2<sup>nd</sup> program  
in definition  
set  
  
 n<sup>th</sup> program  
in definition  
set

A definition, then, is a collection of programs (in the most usual case just one) which depends on a common set of external quantities and which are completely independent with respect to their private internal symbols. The definition set has meaning only at compilation; the independent programs may be dynamically interconnected, among themselves or with programs compiled at another time, in any meaningful way at the time they are executed.

Typing of the definition set is begun by the sequence 'cr tab uc DEFINE'. This first 'DEFINE' insures that the compiler does not retain any symbols mentioned by another user of the system. Each line of a program should be begun with a case punch (uc or lc) and is ended by a carriage return (cr). If a statement is so long that it needs to be broken in typing, the sequence 'cr tab tab tab' provides continuation of the statement onto the next line. 'PROGi' designates a program name. 'PARAMi' designates the parameters of the program, a non-empty list of names separated by commas. The operation '.\_=' followed by the symbol 'SEQ' signals initiation of code generation for the program. Recursive code will be generated (so that a program may use itself) if 'RSEQ' is used instead of 'SEQ'. 'END', typed at the left hand margin and followed immediately by a 'cr', terminates the program, initiates final compiler output of the program, and causes the symbol table limit to be backed up so that the compiler retains only its vocabulary symbols and the external variables of the definition set. The second 'DEFINE' terminates the definition set and causes the symbol table limit to be backed up so that the compiler retains only its vocabulary symbols; all external variables backed over are printed out. 'LEAVE', typed at the left hand margin and followed immediately by 'cr cr', causes exit from the system.

## NAMES

Private names, those invented by a user of the Genie compiler, are formed by the following rules:

- 1) a single lower case Roman letter;
- or 2) an upper case Roman letter, followed by upper case Roman letters, followed by lower case Roman letters, followed by numerals (no embedded spaces).

By rule 1) the following are examples of names:

a i p x

By rule 2) the following are examples of names:

A CAT Fn DDxy I2 PQ29 Dog3

Concatenation of names implies multiplication of the variables specified. The following are not names:

ab A B38 Pt4p M5ef w10

They are interpreted respectively as:

axb AxB38 Pt4xp M5xexf wx10

Any number of characters may be used in a name, but only five are retained by the compiler. If lower case Roman letters are embedded in a name, the first is tallied as two characters.

The names

m Man

are stored as

.M M.AN

Names in the vocabulary of the compiler may not be used by the coder as private names. These include

names of library items -- COL, SIN, LINCT, etc.

names of various machine registers -- B1, CC, T4, etc.

names with special meaning in the Genie language -- as DATA, TRUE, LEAVE, etc.

In alphabetical order, vocabulary names are:

ACOSH	CCSH	CSOLN	FUNCT	MINSE	REPEA
ACCEPT	CDET	CSNH	FXEXP	MITIM	RESUL
ARRAY	CDIV	CSQR	GAMMA	MMPY	ROW
ASIN	CEXP	CSUB	GO	MOD	RTRAN
ASINH	CFEXP	CTAN	IM	MODUL	SCALA
ATAKE	CHISQ	CTNH	INFER	MPATC	SCRIB
ATAN	CINV	CTRAN	INPUT	MPOLA	SET
ATANH	CLENG	CVSPA	INTEG	MPOWE	SIN
B1	CLOG	CXEXP	INV	MRE	SINH
B2	CMADD	DATA	ITIME	MSPAC	SL
B3	CMCON	DEFIN	ITRAN	MSUB	SMDIV
B4	CMCPY	DET	LEAVE	MTAKE	SMMPY
B5	CMPY	DIAG	LENGT	NEO	SOLN
BCD	CMPY	DISPL	LET	NUMBE	SQR
BOOLE	CMSPA	DPUNCH	LGAMM	ODD	STNDV
CADD	CMSUB	END	LINCT	ORTHO	T4
CARTN	CMTAK	EOV	LOG	OUTPU	T5
CACSH	COL	ERASE	LOG10	PAGCT	T6
CALL	COMPL	EVEN	MADD	PAGE	T7
CASN	CONJ	EXECU	MATRI	PLOT	TAN
CASNH	CONTR	EXP	MAX	POLAR	TANH
CATAKE	CONVL	FALSE	MCART	PRESC	TITLE
CATN	COS	FFT	MCMPPL	PRINT	TRAN
CATNH	COSH	FFTC	MCONJ	PUNCH	TRUE
CC	COT	FIX	MCOPY	QCONF	TTAKE
CCEXP	CRCOR	FLEX	MEAN	RANDM	VECTO
CCOL	CROW	FLOAT	MFLT	RE	VREV
CCONT	CSIN	FOR	MIM	READ	VSPAC
CCOS	CSMDV	FORMA	MIN	REAL	Z
CCOT	CSMMP	FTRAN	MINDE	REM	

The following names may be used as private symbols in Genie language but have special meaning in the assembly language:

B6    I    PF    R    S    U    X

Four character strings which are not names have special meaning to the compiler:

and    if    or    not



## NUMBERS

A string of decimal numerals

$$DDD < 2^{14}$$

is an integer. A string of decimal numerals containing either a decimal point '.' or a power point '\*' is a floating point number. The form of a floating point number is illustrated by

$$A.B * C$$

which is interpreted as

$$A.B \times 10^C$$

There may be as many as 14 numerals in A and B combined. C is an integer between -70 and 70; if C is not preceded by a minus sign, it is taken to be positive. Minus signs may precede decimal numbers, integer or floating point, with the usual arithmetic meaning.

A string of 18 or fewer octal numerals immediately preceded by a unary '+'

$$+\phi\phi\phi$$

is a right-adjusted octal configuration. [A '+' between two numbers is binary and does not cause the number which follows it to be octal.]

The following numbers will be understood as shown:

3	decimal, integer
-3.0	decimal, floating point
3.	decimal, floating point
3*8	decimal, floating point
3.0*-8	decimal, floating point
-0.3	decimal, floating point
.3	decimal, floating point
+30	octal

Spaces may be embedded in numbers; they are ignored. Therefore, fields within a number may be separated by spaces for ease of reading. For example, if \_ represents a 'space' punch,

$$3\_641\_209.4\_*\_-8$$

is exactly equivalent to

$$3641209.4 * -8$$

and

+00200\_0130\_0004\_00257

is exactly equivalent to

+002000130000400257

## VARIABLES

In any program, each variable falls into one of three categories: internal, external, or parameters.

Internal variables must be scalars (integer, real floating point, complex, or Boolean), and these are assigned storage within the program. The names of internal variables are not retained outside the compilation of a single program; hence, the same name may be used in more than one program with a different meaning in each of the programs. Labels on statements are also internal variables.

External variables may be either scalar (integer, real floating point, complex, or Boolean) or non-scalar (program, vector, matrix, or array), and all non-scalars must be external. All external variables of a program must appear in the definition set containing that program before any '.\*'. External variables of any one program are the common property of all programs in which they are declared external that are in the machine at running time. The names must have unique meaning throughout the system. During program execution, each external variable has its name on the symbol table (ST, \*113) and its scalar value or non-scalar codeword in the corresponding value table (VT, \*122) entry.

Parameters may be either scalar or non-scalar. If they are non-scalar they must be so declared within the definition set containing the program before any '.\*'. Parameters are neither internal nor external with respect to the program in which they appear, but while running the arguments will fall into one of these categories with respect to dynamically higher level programs. Parameters of a program are only representative of the arguments which will be specified to the program by the dynamically higher level program which uses it while running. Within a system of programs the dynamically highest level program receives control from the operating system and cannot have arguments provided by the system; hence, the dynamically top level program should have one purely dummy parameter, a name that is never referred to in the program. The names of parameters are not retained outside the compilation

## VARIABLES

2

of a definition set; the same name may be used as a parameter for more than one program in a definition set, but for no other purpose in the definition set.

## DECLARATIONS

Declarations are used to describe variables that names represent. The simple form of declaration is illustrated by:

```
VECTOR A
VECTOR A, B, C
VECTORS A, B, C
```

```
|cr      |1st tab
```

A more general form is illustrated by:

```
INTEGER VECTOR A, B, C
INTEGERS VECTORS A, B, C
```

```
|cr      |1st tab
```

One or more declaration words (either singular or plural) are followed by one or more variable names separated by commas.

A variable used in the Genie language is completely described by its:

type	integer, real, complex, or Boolean
shape	Scalar, Vector, Matrix, or Array
and mode	function or not

A scalar is described by a type declaration:

INTEGER or INTEGERS	for integer
REAL or SCALAR or SCALARS	for real floating point
COMPLEX	for complex (Cartesian form)
BOOLEAN	for Boolean

A non-scalar is described by a shape declaration:

VECTOR or VECTORS	for vector	} whose elements are scalars
MATRIX or MATRICES	for matrix	
ARRAY or ARRAYS	for non-scalars whose elements are non-scalars	

and a type declaration which applies to its elements.

A function is described the mode declaration:

FUNCTION or FUNCTIONS	for a private program name
-----------------------	----------------------------

and type and shape declarations which appropriately describe its implicit result, if it has one. Note: Library programs are known to Genie, and need no declarations.

Not all variables need be described by declarations. When

a variable appears on the right side of an equation in the Genie language, its type, shape, and mode will be inferred if they have not been declared:

```

type      real floating point
shape     scalar
mode      non-function

```

The INFER declaration may be used to cause other type and shape inferences:

```

INFER { INTEGER
      { REAL
      { SCALAR
      { COMPLEX
      { BOOLEAN }
      { VECTOR
      { MATRIX
      { ARRAY }

```

where either a type or a shape is given, or both in either order. The range of effect of an INFER is to an INFER which respecifies what it specifies, but not outside a definition set.

The name of every external variable must appear in at least one declaration before any '.\*'. All declarations pertaining to parameters must appear before any '.\*', but they need not appear in any declaration if inference will give a proper description. Declarations pertaining to internal variables must appear within the program to which they belong, and only the type declarations are applicable since all internal variables are scalars.

Not more than one declaration in each group may be applied to a single variable, and not more than one declaration in each group may appear in a single declaration statement.

Thus,

```

      BOOLEAN MATRIX FUNCTION F

```

```

| cr      | 1st tab
is a legal statement, but

```

```

      INTEGER BOOLEAN B

```

```

| cr      | 1st tab
is not.

```

## FUNCTIONS

A function is a program which may be referred to in the Genie language, either for implicit execution as 'F' in the command

$$y=a+F(P)+b$$

or for explicit execution as 'G' in the command

EXECUTE G(Q)

Implicit execution is meaningful only if the function is single valued; in this case its output is not specified in the parameter list. In all other instances explicit execution is required.

The last executed command of a function to be used implicitly must define the output as follows:

RESULT=scalar or non-scalar arithmetic expression

| cr            | 1st tab

In the definition of a function, its parameters are given as an ordered list of those quantities which are supplied as arguments by the program which causes it to be executed. An argument for a parameter which designates a quantity to be calculated by the function must be specified as a simple variable name; other arguments may be given by any arithmetic expression. For example, if  $F(A,B,C)$  is defined such that parameters A and B are used in the calculation of parameter C by the function F, a proper use of F would be  $F(3m^2+n, V_a, P)$ . But  $F(\text{SIZE}, \text{SPAN}, q^2)$  is incorrect since the third argument may not be an expression. Care must be taken that parameters in the definition of a Genie program and arguments in the use of it by other Genie programs are always listed in the same order and agree in number and type.

A function may be sufficiently simple to be defined in one statement. This is done before any '=' and is illustrated by the definition of f in the statement

$$f(x,y)=3ax+a^2y, \quad a=2+x$$

| cr

The function f may then be executed implicitly within the command

sequence of a program,

$$h = k^2 f(m, n)$$

where the closed subroutine  $f$  will be applied to the arguments  $m$  and  $n$ . During compilation, output for  $f$  will be produced independent of that for the other programs in the definition set.

Every Genie program is a function. It may be used as such by any other Genie program. A Genie program begun with 'RSEQ' is a recursive function, one which may use itself. For example, the function FACTL may be executed from within the command sequence for FACTL:

```

      :
FACTL(k) .= RSEQ
      :
      m = FACTL(n-1)
      :
END

```

A recursive function may be executed either implicitly or explicitly, as appropriate to its definition. Genie programs begun with 'SEQ' and functions defined in one statement do not cause recursive code to be generated; they may not use themselves.

All functions except those in the library must be declared in function declarations. If a function is to be executed implicitly and its result is not to be inferred, then its name must appear in declarations to describe the result as well as in a FUNCTION declaration. Thus, the function with its arguments is an operand which must be assigned the type and shape of its output if it is to appear within an arithmetic expression.

A function name is not followed by arguments in a declaration. To specify execution, a function must be followed by arguments, as SIN X2 or CALC(q) or MAP(a+b,VAR). A function name, without arguments, may be supplied as an argument to a function which will do the execution. Thus, the program P may be defined as P(...,F,...), where the parameter F is a function, and call for execution of F(...);



then P may be executed with argument g as  $P(\dots, g, \dots)$  and the result will be execution of  $g(\dots)$  while running.

Note: One inconvenience is associated with this notation. If F1 is a function of a single parameter F2 which is a function, the expression

$\dots F1(F2) \dots$

will be misinterpreted by the compiler. One extra pair of parentheses is required, as

$\dots (F1(F2)) \dots$

if a single parameter is a function.

## CONSTANTS

Internal variables which are constants may be numerically specified by a LET statement within the program. The statement must be given before the name of the constant is used in the commands of the calculation. The form of this statement is illustrated by:

```
LET PI=3.14159
```

```
| cr      | 1st tab
```

This is a message to the compiler which causes the number 3.14159 to be used in the program each time the internal variable name 'PI' is used. A LET statement causes no code to be generated.

The above shows specification of a real floating point value. The variable PI takes on real floating point type.

An integer value may be specified, as

```
LET K=3
```

The variable K takes on integer type.

A complex value may be specified, as

```
LET CVAL=-3.2+/5.19
```

or

```
LET POLE= 1+/0
```

The variables CVAL and POLE take on complex type.

A Boolean value (TRUE or FALSE) may be specified, as

```
LET t=TRUE
```

or

```
LET No=FALSE
```

The variables t and No take on Boolean type.

An octal configuration (right-adjusted) may be specified, as

```
LET MASK=+777777077
```

The + inflection concatenated immediately to the left of a number denotes octal interpretation of the number. The variable MASK should not be used in the Genie language.

A fixed address or codeword address may be specified, as

```
LET #TIME=+200
```

must be used for every numbered scalar, program, vector, or matrix. A Genie program may assign its own name a numerical equivalent, and the tape produced by the compiler will load with codeword at the address specified.

The values of non-scalars may not be specified in a LET statement.

The LET statement may also be used to specify the equivalence of two names. For example

LET ALPHA = BETA

causes 'BETA' to be substituted for 'ALPHA' throughout the program. Similarly

LET COUNT = B5

causes the index register B5 to be used for 'COUNT'.

More than one constant may be specified in a LET statement, if they are separated by commas, as

LET A=3, z=5.41\*-6, #PROG=+127, TIME1=TIME2

There are three other commands which identify names with values. They are explained later: BCD, NUMBERS, and FORMAT in the section on data commands. These commands are non-executable and must be transferred around, and must therefore be used with care.

## REMARKS

Printed comments in compilation output listings may be obtained by using the REM statement within the program, as illustrated by

```
REM COMPUTE_FIRST_VALUE
```

or

```
REM
```

```
COMPUTE_FIRST_VALUE
```

```
| cr      | 1st tab
```

where \_ indicates a space typed within the remark. 'REM' is followed immediately by a single space or 'cr' which is not part of the remark, and then all following characters are taken as remark text. The statement may be continued to succeeding lines at the 3rd tab position by using the 'cr tab tab tab' sequence. The form of REM in which the text begins at the left margin causes remarks to stand out more vividly on program listings.

The REM statement does not introduce any data into the final program; its only effect is to cause the remark to be printed in the compilation output listing.

## COMMAND SEQUENCE

All statements of a program from the `'.'` to and including the `'END'`, except `'LET's`, remarks, and declarations, cause code to be generated. Such statements are called commands. The occurrence of a label on a command causes a command sequence to be initiated. The ordered set of all command sequences of the program is called the command sequence for the calculation. Each command falls into one of four categories; arithmetic, control, input-output, or data. These will be discussed in separate sections.

Any command may be labelled. The label is typed at the left-hand margin, as `'CALC'` in the command

```
CALC      A=B2+B+3.2,  B=W+5.1  
| cr      | 1st tab
```

## ARITHMETIC EXPRESSIONS

The righthand side of an equation in the Genie language must be an arithmetic expression. An arithmetic expression is just an operand. A scalar constant, a variable, an inflected variable, or a function name followed by a parenthesized list of arguments is an operand. [A single argument given as a simple variable name need not be enclosed in parentheses.] A pair of operands joined by an operation (where the triplet left operand, operation, right operand is defined in Genie) is an operand.

Any operand may be enclosed in parentheses to dictate order of computation within an expression in the conventional manner. Order is also implied by relative rank of operations. In order of decreasing rank, i.e., the most binding first, the arithmetic operations are:

unary inflections: -, |...|, and 'not'  
subscription  
exponentiation  
x and /  
+ and binary -  
relations: =, ≠, <, ≤, ≥

The triplets of operands joined by an arithmetic operation which are permitted in an arithmetic expression on the righthand side of an equation are given in the following paragraphs.

- 1) +, -, x, / between integer, real floating point, or complex scalar operands.

If the operands are both integer or both floating point or both complex, the result is of the same type. If one operand is complex and the other is not, the non-complex operand is made complex before the operation is carried out, and the result is complex. If one operand is floating point and the other is integer, the integer is floated before the operation is carried out, and the result is floating point.

- 2)  $+/$  between integer or real floating point scalar or non-scalar operands.

This is the explicit representation of a complex quantity in Cartesian form, as  $x+/y$  (written  $x+iy$  in mathematical notation). The result is complex with real and imaginary parts real floating point. The shape (scalar, vector, or matrix) of the parts determine the shape of the result; both parts must be of the same shape, and non-scalars must have the same dimensions. If the operands joined by  $+/$  are expressions, they must be enclosed in parentheses. If the operand  $x+/y$  is combined arithmetically with other terms, it must be enclosed in parentheses.

$-/$  between integer or real floating point scalar operands.

The complex scalar  $x-/y$  is simply  $x+/-y$ .

- 3)  $+(or)$ ,  $-(symmetric\ difference)$ ,  $\times(and)$ ,  $/(symmetric\ sum)$  between Boolean scalar operands.

Combination of Boolean operands yields a Boolean result, by the following rules:

- $+$  FALSE if both operands FALSE, otherwise TRUE
- $-$  TRUE if operands differ, FALSE if operands the same
- $\times$  TRUE if both operands TRUE, otherwise FALSE
- $/$  TRUE if operands the same, FALSE if operands differ

The octal representations for the Boolean values are

TRUE      0077777777777777

FALSE     0077777777777776

- 4)  $+$ ,  $-$ ,  $\times$  between non-scalar operands containing integer, real floating point or complex elements.

Standard conventions apply as to restrictions on dimensional compatibility, and the operands must be in standard form.\* Addition or subtraction of two vectors or two matrices yields a vector or a matrix respectively. Multiplication of matrices yields a matrix. Multiplication of vectors yields the scalar product which is

a scalar. Multiplication of a vector and matrix yields a vector. If the operands are of the same type, the result is of that type. If the operands are of different types and one is complex, the result is complex. If one operand is integer and the other floating point, the result is floating point.

- 5)  $\times$  between integer, real floating point, or complex scalar and integer, real floating point, or complex non-scalar.

The scalar may be on the left or the right of the non-scalar, which must be in standard form.\* The result has the same form as the non-scalar operand, vector or matrix. If the operands are both integer or both floating point or both complex, the result is of the same type. An integer operand is floated before combination with a floating point operand, and the result is floating point. An integer operand is floated and then made complex before combination with a complex operand, and the result is complex. A floating point operand is made complex before combination with a complex operand, and the result is complex.

- 6) Division of an integer, real floating point, or complex non-scalar by an integer, real floating point, or complex scalar.

The non-scalar must be in standard form.\* The result has the same form as the non-scalar operand, vector or matrix. If the operands are both integer or both floating point or both complex, the result is of the same type. An integer operand is floated before combination with a floating point operand, and the result is floating point. An integer operand is floated and then made complex before combination with a complex operand, and the result is complex. A floating point operand is made complex before combination



with a complex operand, and the result is complex.

- 7) Implied multiplication between operands which appear immediately next to one another, not separated by an operation. The same rules apply as for the explicit  $\times$ .
- 8) Exponentiation between integer, real floating point or complex scalar operands.

If either or both operands are complex, the result is complex. If neither operand is complex but either or both operands are floating point, the result is floating point and the base may not have a negative value. If both operands are integers, the result is an integer, zero if the base is  $> 1$  in absolute value and the exponent has a negative value. Note that  $A^B$  is typed 'A sup B sub', using the superscript and subscript keys on the flexowriter. The counter associated with these carriage moving keys should be set to zero before starting a program and must return to zero before the cr which ends each command.

- 9) Exponentiation of a short logical operand by an integer.

Short logical words are 15-bit configurations whose bits are numbered 1 to 15 from left to right. In particular SL (the sense light register) is in the vocabulary of the compiler and falls into this category. The result of exponentiation of such an operand by an integer, as  $SL^k$ , is Boolean, TRUE if bit  $k$  of SL is 1 (on) and FALSE if it is 0 (off). The value of the bit addressed is not affected by the operation. The user may also exponentiate a private variable which has been declared BOOLEAN.

- 10) Exponentiation of a square integer or real floating point matrix to an integer power.

If the matrix is integer it will be floating before exponentiation. The matrix must be in standard form.\* The result is always a floating point matrix. If P

is the power and  $P < 0$ , the inverse is computed. If  $|P| > 0$ , multiplication occurs  $|P-1|$  times. If  $P=0$ , the result is the unit matrix.

- 11) Subscripting of a vector by an integer scalar operand or of a matrix by a pair of integer scalar operands separated by commas.

The result is an element of the vector or matrix and is of the same type (integer, real floating point, complex, or Boolean) as the non-scalar of which it is an element. The expression  $A_B$  is typed 'A sub B sup' and return to zero carriage level must be observed as for exponentiation.

- 12) Any non-scalar may be subscripted with a total of five integer subscripts separated by commas. The operand is indirectly addressed after  $B_1, \dots, B_5$  are loaded with the subscripts. An Array may be subscripted at both levels in one expression, e.g.  $\dots(A_{I,J,K})_{L,M}\dots$ , where A in an Array, is a reference to element L,M of the matrix  $A_{I,J,K}$ . The placement of the parentheses indicates the break point in the structure and the subscripting procedure is restarted with  $B_1$ . The parentheses are not necessary for the first level, e.g.  $\dots B_{K,L}\dots$ , where B is an Array, is a reference to non-scalar  $B_{K,L}$ .

- 13) Relations  $=, \neq, <, \leq, \geq$  between integer or real floating point scalar operands.

Combination of integer or floating point operands with a relational operator yields a Boolean result, TRUE if the two operands stand in the specified relation to each other, FALSE otherwise. If the operands are not both integer or both floating point, the integer operand is floated before the comparison is made. If r and r' are relations, the form  $ArBr'C$  is tempting but not permitted; an equivalent form is  $(ArB) \times (Br'C)$ . A precise sequence of typed characters

is required:

$\neq$  is typed ' = backspace uc | '

$<$  is typed ' < backspace uc | '

$\leq$  is typed '  $\leq$  backspace | '

Note that the relations  $>$  and  $\geq$  are not available, but  $>$  is equivalent to  $\neq$  and  $\geq$  is equivalent to  $<$ .

- 14) Unary - applied to an integer, real floating point, or complex scalar operand.

The negation of the operand takes place before it is combined with any other across a binary operation, except exponentiation and subscription.

- 15) Absolute value of an integer or real floating point scalar operand.

This inflection is denoted by absolute value bars '|'| before and after the operand. These bars are simply parentheses that cause the quantity inside to be taken with positive sign.

- 16) Unary 'not' or - applied to a Boolean scalar operand.

The complementation of the Boolean operand takes place before it is combined with any other across a binary operation, except exponentiation and subscription. If the Boolean scalar has the value TRUE, then not A has the value FALSE; if A has the value FALSE, not A has the value TRUE.

\* The standard form for vectors and matrices is that handled by VSPACE, MSPACE, and the Genie input-output commands. Generation and input-output of non-standard forms can only be handled by explicit use of SPIREL facilities.

## ARITHMETIC COMMANDS

The form of a simple arithmetic command is illustrated by:

A=arithmetic expression

| cr        | 1st tab

The form of a compound arithmetic command is illustrated by:

A=arithmetic expression, B=arithmetic expression, ...

| cr        | 1st tab

where more than one equation appears in the command.

If there are no interdependencies among the equations of a command, the equations are coded by Genie in the order given. If there are interdependencies, the first equation will be coded last and preference will be given to coding the remaining equations from right to left; for the second and any following equations, if the  $i^{\text{th}}$  depends on the  $j^{\text{th}}$  and  $i > j$  (counting from left to right), then the  $j^{\text{th}}$  equation will be coded before the  $i^{\text{th}}$ . So the second and following equations may well be used to define subexpressions of the first (or primary) equation, producing code that will run more efficiently and copy that will be more readable. An example in which reordering will take place is

y=a+b, a=5c/d, b=6, c=b+4

| cr        | 1st tab

The code generated will evaluate b, then c, then a, then y. On the other hand, the equations in

M=P+Q, a=3, i=j+1

are not dependent upon each other and will be coded in the order given.

The variable on the lefthand side of an equation may be a scalar, or a non-scalar, or a subscripted non-scalar (denoting a scalar element of a vector or matrix). All lefthand side variables in a command must be distinct, no scalar or non-scalar defined more than once. More than one element of the same non-scalar may be defined in one command.

The '=' joining lefthand side to righthand side of an equation causes storage of the computed righthand side into the loca-

tion or array specified on the lefthand side. Compatibility of types is checked for at time of compilation, and an error message is printed out if incompatibility of the two sides is detected. In every case the righthand side dominates and will be stored as calculated, no conversion taking place. If the righthand side is non-scalar, the storage addressed by the codeword on the lefthand side is freed before the store across the '=' takes place.

Genie has the ability to apply the commutative laws of arithmetic to reorder the terms of an expression to provide calculation using a minimum number of temporary stores. In the coding for a non-complex scalar expression, the compiler may use the T-registers of the computer for temporary storage. Push-down storage addressed by index register B6 is also used for this purpose. When profitable, the T-registers are used by the compiler for non-complex scalar variables that are referred to often in an equation. The codeword at machine address 10 (octal) is used in the code generated by the compiler as an accumulator for real vectors and matrices produced in the course of evaluating the righthand side of a non-scalar equation. This address may not be used by a coder. The accumulator for complex non-scalars is named CSTAR. Temporary storage for non-scalars is always on the B6-list.

## CONDITIONAL ARITHMETIC COMMANDS

A simple arithmetic command may be of conditional form, as illustrated by

$$A = E_1 \text{ if } P_1, E_2 \text{ if } P_2, \dots, E_n \text{ if } P_n, E_{n+1}$$

| cr            | 1st tab

where each  $E_i$  is an arithmetic expression and each  $P_i$  is a predicate which is either true or false. The code that is generated will evaluate  $A$  as  $E_i$  for the least  $i$  for which  $P_i$  is true. If every  $P_i$  is false, then  $A$  is evaluated as  $E_{n+1}$ . If  $E_{n+1}$  is omitted, then  $A$  is not evaluated at all if every  $P_i$  is false.

Boolean valued expressions are predicates, as in the following examples:

$$K = 1.0 \text{ if } B \leq C, 2.0 \text{ if } x < -12.9, 3.0$$

$$K = 1.0 \text{ if not } (SL^n), 3.0$$

$$K = 1.0 \text{ if } SL^5 + \text{not } (SL^n)$$

| cr            | 1st tab

Boolean valued expressions joined by the operations 'and' and 'or' form predicates, as in the following example:

$$K = 1.0 \text{ if } (B \leq C \text{ or } |C + D| \neq 3.72) \text{ and } SL^5 + \text{not } (SL^n), \\ 2.0 \text{ if } x < -12.9, 3.0$$

| cr            | 1st tab | 2nd tab | 3rd tab

The most binding first, the operations are ordered as follows:

arithmetic operations

'and'

'or'

Parentheses may be used, as in the above example, to dictate computational order.

The predicate form  $F_1 \text{ r } F_2 \text{ r' } F_3$  is tempting but not permitted. An equivalent permissible form is

$$F_1 \text{ r } F_2 \text{ and } F_2 \text{ r' } F_3$$

$$\text{or } (F_1 \text{ r } F_2) \times (F_2 \text{ r' } F_3)$$

Two exceptional Boolean predicates are 'EOV', asking if the exponent overflow light is on, and its negation 'NEO'; neither of these may be inflected by 'not'. Both of these tests turn the

## CONDITIONAL ARITHMETIC COMMANDS

2

light in the indicator register off.

A conditional arithmetic equation must stand alone as a command. It may not be grouped with other equations in a compound arithmetic command.

## TRANSFER CONTROL COMMANDS

Code is generated so that the commands of the program are normally executed in the order written. An explicit variation in this order is indicated by a transfer command, illustrated by

```
CC = #LOOP or GO TO LOOP
```

```
| cr      | 1st tab
```

Here 'CC' is the mnemonic for the control counter which is normally stepped sequentially through the orders of the code. 'LOOP' is a label on a command of the program, the command to which control will be passed by this transfer command. Note that 'END' is a label in every program and may be transferred to for exit from the program. The inflection '#' is required in this context to indicate that the address corresponding to LOOP, and not the contents of the location whose address is LOOP, is to be calculated on the righthand side. The '#' inflection is analagous to the 'a' bit in APl.

The conditional transfer command provides variation in the order of command execution depending upon the truth values of predicates. The form of this type of control command is shown by

```
CC = #A1 if P1, #A2 if P2, ..., #An if Pn, #An+1 or
```

```
| cr      | 1st tab      GO TO A1 if ...etc.
```

where the A<sub>i</sub> are labels within the program and the P<sub>i</sub> are predicates. The code generated causes CC to be evaluated as the first #A<sub>i</sub> for which P<sub>i</sub> is true. If no P<sub>i</sub>, for i=1, 2, ..., n, is true, CC is evaluated as #A<sub>n+1</sub>. The term #A<sub>n+1</sub> may be omitted from the command, in which case CC is unchanged if all P<sub>i</sub> are false, so that no transfer is made. The predicates P<sub>i</sub> are of the form described in the section on conditional arithmetic commands.



## LOOP CONTROL COMMANDS

Loops may be realized in Genie language by a combination of arithmetic commands and transfer control commands. A concise notation for a popular loop structure is provided by the loop control commands. The commands of a loop are parenthesized by the FOR and REPEAT commands of the form

```
FOR P=A, B, C
    commands of the loop
```

```
REPEAT
```

```
| cr      | 1st tab
```

The elements of the FOR command are  
parameter of the iteration, P  
initial value, A  
increment, B  
final value, C

All elements must be scalars, either integer or floating point. In execution, the loop is traversed for  $P = A + kB$ , for all  $k = 0, 1, 2, \dots$  such that

$$P \leq C \text{ if } B > 0$$
$$P \geq C \text{ if } B < 0$$

The element P must be given as a simple variable name. The elements A, B, and C may be given as constants or arithmetic expressions of integer or floating point type. Only if B and C are given as simple variable names may their values change during execution of the loop. Otherwise, B and C retain their values on entry to the loop throughout the execution of the loop. For example, in the loop

```
FOR COUNT = FIRST, M+N, LAST
```

```
  :
```

```
  N=A+B
```

```
  :
```

```
REPEAT
```

the increment value will remain constant, as computed on entry to the loop.

In the REPEAT command, 'REPEAT' is followed immediately by a

'cr'. A REPEAT must be written for every FOR.

If addressed from outside the loop, the iteration parameter has the value it had upon exit from the loop.

Loops may be nested to any level, but distinct iteration parameters must be used at each level within a nest. The 'REPEAT' is considered to be within the loop which it terminates; the 'FOR' is not. Transfer of control may be made from a command within a loop to another command within the loop or to a command outside the loop. Transfer from outside a loop to the FOR command is permitted, but transfer from outside a loop to a command within a loop is not permitted.

Any 'FOR' or 'REPEAT' may be labelled for purpose of transfer to it. The compiler generates the label ' $\leftarrow$ FORn' on each FOR command and ' $\leftarrow$ RPTn' on the corresponding REPEAT command,  $n = 1, 2, \dots, 9, a, b, \dots$  in each program. A coder's label will be used instead if it appears. Thus, FOR and REPEAT commands begin command sequences whether or not they are labelled by the coder.

The machine index registers B3, B4, B5 may be used as iteration parameters in loops and will cause significantly more efficient code to be generated, especially when a constant increment  $= \pm 1$  is specified. The section on fast registers discusses coder usage of machine registers.

## STORAGE CONTROL COMMANDS

Before a vector or matrix is referred to dynamically in a program it must be created, either initially from paper tape or dynamically while running.

In a Genie program, to create, or take space for, the vector named VNAME of length NELTS elements the following command is used:

```
EXECUTE VSPACE(VNAME, NELTS)
```

```
| cr      | 1st tab
```

The vector VNAME contains zeroes initially. To create, or take space for, the matrix named MNAME of NROWS rows and NCOLS columns the following command is used:

```
EXECUTE MSPACE(MNAME, NROWS, NCOLS)
```

```
| cr      | 1st tab
```

The matrix MNAME contains zeroes initially. The dimension arguments in both commands are integers.

The dimension arguments may be computed dynamically, so that sizes of vectors and matrices may vary from run to run. In fact, the dimension of an array may vary during a run by use of a creation command to 'recreate' an array which already exists; the old copy is automatically erased before the new one is formed.

To explicitly erase, or free the space occupied by, a vector or matrix named ARRAY on which the calculation no longer depends the following command is used:

```
ERASE ARRAY
```

```
| cr      | 1st tab
```

Also a single ERASE command may be applied to more than one non-scalar, as illustrated by:

```
ERASE VNAME, MNAME, ARRAY
```

```
| cr      | 1st tab
```

The erasure of a vector or matrix causes the storage occupied to be returned to a common pool, that from which storage is obtained for the creation of vectors and matrices. This pool is managed by STEX, the storage exchange program in SPIREL (explained in detail

in the literature on SPIREL), and it is called the STEX domain. STEX may move items within its domain to concentrate space if necessary to satisfy requests for space.

## EXECUTE CONTROL COMMANDS

The command

```
EXECUTE PROG(PARAM)
```

```
| cr      | 1st tab
```

causes control to be transferred to the program whose name is denoted by 'PROG' in this illustration. 'PROG' must have been declared as a function outside the command sequence for the calculation. 'PARAM' denotes a list of one or more parameters separated by commas. Parameters may be arithmetic expressions unless they designate quantities which are to be calculated by the function, in which case they must be simple variable names. Control is returned from PROG to the next command in the sequence. The interpretation given to the EXECUTE command by Genie is parallel to that for the arithmetic command, the information to the right of the space after the EXECUTE corresponding to that after the first '=' in an arithmetic command. Thus, a simple conditional EXECUTE command is allowed, such as

```
EXECUTE A(P) if a < b + c, B(Q)
```

```
| cr      | 1st tab
```

And a compound unconditional EXECUTE command is allowed, such as

```
EXECUTE SUM(x,y), x = 2a/b, y = ab, b = 4
```

```
| cr      | 1st tab
```

## INPUT-OUTPUT COMMANDS

The input-output commands are:

DATA list	READ list	†PAGE list
PRINT list	INPUT list	ACCEPT list
PUNCH list	OUTPUT list	TITLE string
DPUNCH list	DISPLAY list	

| cr        | 1st tab

where 'list' denotes a collection of names separated by commas. Any name may be that of a scalar, other than fast registers, or of a standard vector or matrix or of a function. Expressions may not appear in the argument list, so vector and matrix elements in the subscript notation may not be designated.

The DATA command provides reading of manually punched signed decimal numbers from paper tape. The name of any type of variable may appear in the list, and any name may have been assigned a machine address in a LET statement. When the paper tape is read, if a decimal point appears the number will be converted to floating point within the machine; the absence of a decimal point causes conversion to integer form. Every number on the tape must be followed by a carriage return, tab, or comma. Integers greater than or equal to  $2^{14}$  in absolute value are meaningless; floating point significance to more than 14 places is not meaningful. A floating point number may be followed by the sequence '\* signed integer' which will cause it to be multiplied by 10 to the signed integer power upon conversion. The magnitude of such numbers must be greater than  $10^{-70}$  but less than  $10^{70}$ . The absence of a sign on a number implies positive sign. Then

punched    328cr	converts to	integer 328
46.9cr		floating point 46.9
.469*2cr		floating point 46.9
-5391cr		integer -5391
-69.*-1cr		floating point -6.9

†Blank or numbers 1 through 7 only

Integers and real floating point scalars are punched as single decimal numbers in the appropriate format; complex scalars are punched as real part followed by imaginary part, both floating point. A vector of length  $n$  is punched as the sequence of  $n+1$  decimal numbers: integer  $n$ , first element, ...,  $n^{\text{th}}$  element. A matrix of  $m$  rows by  $n$  columns is punched as the sequence of  $mn+2$  numbers: integer  $m$ , integer  $n$ , element (1,1), element (1,2), ..., element (1, $n$ ), element (2,1), ..., element (2, $n$ ), ..., element ( $m$ ,1), ..., element ( $m$ , $n$ ). When the DATA command is executed, the proper tape is assumed to be in the reader. If sense light 14 is off, the line

DATA NAME

| cr | 1st tab

will be printed out for each quantity read, where 'NAME' is as designated in the program containing the READ command. Thus, printer monitoring of DATA applied to parameters bears the dummy parameter name, not the name of the argument supplied as the parameter.

The PRINT command provides decimal output on the fast line printer of any named scalar or non-scalar quantities. These are labelled by the name given in the argument list. Any name may have been assigned a machine address in a LET statement. Scalars are printed four per line. Vectors are printed five elements per line, the leading element index in octal at the left of each line. Matrices are printed by row, five elements per line, the leading column index in octal at the left of each line. Complex variables are printed as real part followed by imaginary part; the name of the variable will be given with the real part, and "ditto" (printed '-----') will label the imaginary part.

The PUNCH command and the READ command may be applied only to external variables and to parameters representing arguments which

at the time of execution are external in some dynamically higher level program fall into this category. Care must be taken to apply these commands properly to parameters as there are no checks built into the compiler or input-output program to insure that scalars internal to some program are not punched or read. A name which has been assigned a machine address in a LET statement may appear in the list for PUNCH or READ. PUNCH provides, for each variable listed, as many control words as are necessary to recreate the form of the variable at a later read time. The content of the variable is punched in hexad with checksum format. These output tapes may be loaded through SPIREL or they may be read with a READ command. The READ command will read any tape produced by PUNCH. Also, READ will read any scalar, standard vector, or standard matrix punched with name by use of SPIREL directly.

The DISPLAY command provides decimal display of named scalars on the storage scope at the console. These are labelled by the name given in the argument list. Any name may have been assigned a machine address in a LET statement. As many as eight lines will be displayed by a single command. Real scalars are displayed one per line. Complex scalars are displayed on two lines, real part with the name of the variable followed by imaginary part with the name "ditto" (written '+++++'). Non-scalars may not be displayed.

The INPUT command and the OUTPUT command provide input and output of named scalars and non-scalars through programs supplied by the user. Any name may have been assigned a machine address by a LET command. For each variable named in an INPUT command control is passed to the program named INPUT; for an OUTPUT command, the program named OUTPUT is used. A complex variable is handled as two variables, the real part with the name of the complex variable and the imaginary part with the name "ditto". Details about the INPUT and OUTPUT programs are given in the library literature.



Formatted printer output may be obtained by use of the command

```
EXECUTE SCRIBE(A1,...,AK,F)
```

```
| cr      | 1st tab
```

where A1,...,AK is the list of arguments to be printed, and F is the name of a FORMAT statement to be used. Any argument in A1,...,AK may be a simple name or an expression. The program SCRIBE is in the library, and its use is fully described in the library literature. A FORMAT statement gives text which will be printed directly by SCRIBE and dummy variables which will be replaced by argument values.

Page control and headings are provided with formatted printer output by use of the command

```
EXECUTE PRESCRIBE(A1,...,AK,F,N,NAME,LIMIT)
```

```
| cr      | 1st tab
```

where A1,...,AK,F are just as for SCRIBE, N is the number of blank lines after SCRIBE output, NAME is the name of a FORMAT statement containing pure text or a vector of BCD data to be used in the heading on each page, and LIMIT is the number of lines per page of output. The program PRESCRIBE is in the library, and its use is fully described in the library literature.

Additional forms of input and output may be obtained by use of SPIREL programs directly, but those provided by the input-output commands should be sufficient for a large number of problems. Also see the TITLE and PAGE commands on p.5.

The DPUNCH command may be applied only to external variables as explained earlier in the PUNCH command. DPUNCH provides standard decimally formatted punched tape to be later read by a DATA command only. Mixed integer and real data may be punched from scalars, vectors or matrices.

The TITLE command allows the printing of a string of literal symbols for labeling pages like SCRIBE only with greater ease. Two examples are given below. One would write:

```
TITLE PRINT ONE LINE HERE
```

```
TITLE
```

```
PRINT ANOTHER LINE HERE ALSO
```

```
| cr      | 1st tab
```

The above would cause the following to be printed:

```
PRINT ONE LINE HERE
```

```
PRINT ANOTHER LINE HERE ALSO
```

```
↑
```

(first printer position)

The PAGE command allows the page to be moved to any position or by any amount easily. The 'list' consists of the integers 1,2,...,7 or no list, i.e., blank. The interpretation of the integers is given by the table below:

integer	→	move	to	next ____ page	
1	→	move	to	next 1/66 page	(one space)
2	→	move	to	next 1/22 page	
3	→	move	to	next 1/11 page	
4	→	move	to	next 1/6 page	
5	→	move	to	next 1/3 page	
6	→	move	to	next 1/2 page	
7	→	move	to	full page	(page restore)

If the list is blank, the page is restored.

The ACCEPT command provides reading of data input through the console typewriter. The name of any type variable may be included in the list. Data may be entered when the blue light on the typewriter comes on; each line is processed before another may be typed. Decimal numbers are handled as in the DATA command; octal numbers must be preceded by a + sign. T... or F... is typed for the Boolean values. All values must be separated by commas and a line is terminated by a carriage return. For a vector or matrix of size n or nxm, n or nxm values must be typed. To change the size of a non-scalar, the new dimension(s) is enclosed by parentheses (n) or (n,m), and followed by the values to be stored. A matrix is typed by rows (as it is read in DATA). For example, where A is a vector 3 long, B is Boolean, and C is scalar:

ACCEPT A,B,C

typewriter input: -234.0, 8.34\*4, .62023, T, +0142000000 cr  
stores the first 3 values in A, -Z in B, and the octal number in C.

typewriter input: (4), 8.0, 9.0, -10.0, 11.0, FALSE, 345 cr  
erases array A and creates a new one of length 4, stores the next 4 values in A, -1 in B, and decimal value 345 in C.

typewriter input: →, T, +002345 cr  
leaves A as it is, stores -Z for B, and the octal value for C. The "forward arrow" is inserted whenever an item in the list is to remain unchanged. If sense light 14 is off, the line

ACCEPT NAME

will be printed out for each quantity in the list (as in the DATA command).

## LIGHT CONTROL COMMAND

The SET command provides program control over sense light setting. It is illustrated by

SET not SL<sup>5</sup>, SL<sup>9</sup>, SL<sup>1</sup>, not SL<sup>15</sup>

|cr |1st tab

Any number of sense lights may be set. The notation 'SL<sup>i</sup>' causes SL<sup>i</sup> to be turned on; 'not SL<sup>i</sup>' causes SL<sup>i</sup> to be turned off. In 'SL<sup>i</sup>' i must be numeric and may range from 1 to 15. The lights are set in the order mentioned.

Data commands cause generation of words in the program which are not instructions. These commands are not executable and all but **FORMAT** must be transferred around.

Alphabetic information for output on the printer may be defined by the BCD command, as illustrated by

```
MESS1      BCD  _ _TEMPUS_FUGIT
```

or

```
MESS1      BCD
_ _TEMPUS_FUGIT
```

```
| cr      | 1st tab
```

where 'BCD' is followed immediately by a single space or a 'cr' which is not part of the data, and \_ indicates a typed space. The command may continue onto succeeding lines at the 3rd tab position by use of the 'cr tab tab tab' sequence. A space is inserted by Genie between the last character of one line and the first of the next line. At the place such a BCD command appears in the command sequence for the program, the printer code for the information is inserted in the code for the program, nine characters per word. The label (if any) on the BCD command is associated with the first word of data.

A block of numeric data may be defined by the NUMBERS command, as illustrated by

```
CONST      NUMBERS 36.5, -2*8, 6, +774777
```

```
| cr      | 1st tab
```

In the program Genie generates, in this case,

floating point 36.5 at CONST

floating point  $-2.0 \times 10^8$  at CONST+1

integer 6 at CONST+2

octal 774777 (right-adjusted) at CONST+3

One or more real numbers (each but the last followed by a comma) are listed; complex numbers may not appear in the list. The list may be extended onto the succeeding lines by use of the 'cr tab tab tab' sequence. The numbers are inserted into the program in the order

given, one per word. The label (if any) on the NUMBERS command is associated with the first word of data.

Formats for the printer output programs SCRIBE and PRESCRIBE are defined by the FORMAT command, as illustrated by

```
LINE      FORMAT ddd  ITERATIONS,  CASE aa,   K=bb,   T=-d.ddce+d  
or
```

```
LINE      FORMAT  
ddd  ITERATIONS,  CASE aa,   K=bb,   T=-d.ddce+d  
| cr          | 1st tab
```

where 'FORMAT' is followed immediately by a single space or a 'cr' which is not part of the data. The label on the FORMAT command is the name of the FORMAT which is an argument to the output programs. The format data is a "dummy line" of printer output; lower case letters and the characters '.', '+', '-' with 'd' form dummy variables for which argument values are substituted when printing; the rest of the format data is text which is printed directly. SCRIBE and PRESCRIBE are programs in the library; their use and the details of format specification are explained fully in the library literature.

## FAST REGISTERS

It is never necessary to use machine registers in the Genie language. But their use is permitted, with certain restrictions and with effect that more efficient code may be obtained.

T7 should never be used in the Genie language.

T6, T5, and T4 may be used as the names of scalar variables within a command. The compiler will not make use of any T-register mentioned by the coder, and code efficiency may be increased by explicit assignment of auxiliary variables to these fast registers. The values in T6, T5, T4 are not preserved by Genie from one command to another as they are subject to use by the compiler in any command in which they are not explicitly mentioned by the user.

The index registers B3, B4, B5 may be used as the names of scalar integers. These are disturbed by Genie-generated code only to address elements of arrays of more than two dimensions. (Non-standard subscripting is discussed in the section on arithmetic expressions.) Efficiency of code is gained if these registers are used as subscripts or as iteration parameters of loops with constant increment  $\pm 1$ . The index registers B1 and B2 may be used only if the user understands Genie coding conventions are explained in another section and can accurately anticipate the use of these registers by Genie generated code. The registers B6 and PF may not be used in Genie language.

## ASSEMBLY LANGUAGE

In a Genie program, instructions in the AP2 assembly language may be interspersed at will with commands in the Genie language. AP2 is discussed in detail in the assembly language literature.

The following names identify fast registers in both Genie language and AP2:

T4	T6	CC	B2	B4
T5	T7	B1	B3	B5

The following names identify private quantities in Genie language and fast registers in AP2

R	B6	U	I
S	PF	X	

Therefore, a private name I in Genie language may not be addressed in AP2 code.

Operations without mnemonics in the AP2 vocabulary may be coded in octal, as

+45061      #15

| cr      | 1st tab | 2nd tab | 3rd tab

Or an operation code mnemonic may be assigned with a LET statement, as

LET #QSR = +45061

Then the instruction

QSR      #15

could be used instead.

In AP2 commands, the coder may make use of the fast registers, taking care to preserve the value of PF for reference to parameters and to use B6 for temporary push-down storage only. Entire functions may be written in the assembly language, but the user must first understand various Genie coding conventions, as discussed in a later section.

Normally for a Genie program initial and terminal program sequences and code to preserve parameter addressing are automatically generated by the compiler. For some programs coded predominantly in AP2, it may be desirable to avoid generation of or-



ders not explicitly coded. This may be accomplished by using 'ORG' in place of 'SEQ', as

```
PROG(PARAM).= ORG
```

```
| cr
```

to start the command sequence for the program. The first instruction of the program will be the first explicitly coded. The only words in the program generated automatically by the compiler are cross-references to external quantities and a one-word END program sequence:

```
END      TRA      Z
```

The programmer must code parameter set-up for the program, maintain PF and B6 by Genie coding conventions.

## PUNCTUATION

Reference to rules of punctuation for use in the punching of Genie programs has been made in other sections. A few generalities and notes here may help the user to avoid some of the most common mistakes.

Every tape must begin with a 'cr' punch and a case punch for proper interpretation.

Every line should begin with a case punch so that it does not depend on the case at termination of the preceding line, and editing of tapes will be thus simplified.

Spaces may appear anywhere but within a name; they will be ignored.

Backspaces are ignored except within the sequence of punches for negated relations.

The superscript and subscript punches should be used only where meaningful; the sequences 'sup sub' and 'sub sup' are not equivalent to no punch at all and will not be accepted by the compiler.

The carriage counter should be set to zero before typing a program and must return to zero before the 'cr' which ends each statement.

A statement is continued onto second and succeeding lines by the sequence of punches 'cr tab tab tab'.

The operation '.\*' must be punched as just those two characters in succession.

The negated relations require specific sequences of punches for proper interpretation:

⊢ is punched '= backspace uc |'  
⊥ is punched '< backspace uc |'  
⊤ is punched '≤ backspace |'

The operations 'not', 'and', 'or', 'if' are punched in lower case and must contain no superfluous punches. All other "words" in the vocabulary of the compiler are punched fully in upper case letters.

Statement labels, the program name, function definitions 'END', and 'LEAVE' are typed at the margin; alternatively, program names and function definitions may be typed at the 1st tab position.

Since 'SEQ', 'END', and 'DEFINE' end statements, they must be followed immediately by a 'cr' punch.

Declaration identifiers, 'DATA', 'EXECUTE', 'FOR', 'LET', 'NUMBERS', 'PRINT', 'PUNCH', 'DPUNCH', 'READ', 'SET' may be followed by either a space or a tab punch.

'BCD', 'FORMAT', 'REM' may be followed by a space, a tab, or a carriage return punch.

For compilation to be terminated properly 'LEAVE' must be followed immediately by two 'cr' punches.

## COMPILEATION PROCEDURE

A Genie program is compiled by exercising option #6 in the PLACER system.

Compilation output on the printer consists of error messages, program listing, and symbol tables. These are discussed below. Compilation provides a punched paper tape to be loaded under SPIREL control. Compilation options are also discussed below.

Error messages. Genie error messages refer to carriage return number on the PLACER listing of the program. During compilation the carriage return number for the line being compiled is displayed in FT (the from-tape register). This can be useful if compilation problems arise with no error message. If a single command, statement, or instruction is continued onto more than one line, the carriage return number for the last line will pertain throughout.

Program listing. Four columns are printed, giving:

- (a) The symbolic location (if any).
- (b) The relative location of the word in the program, in octal.
- (c) The instruction in octal, broken into fields, with tag.
- (d) The symbolic address (if any).

Cross reference words and internal storage are listed after the instructions of the program, one per word with name, relative location, and content for each. The variables referenced relative to PF are then listed with name and PF increment.

Symbol tables. For each program a symbol table of internal names is printed. Of interest are columns which give the name and the relative location in the program (two to the right of the name). The column to the right of name contains descriptive information about the variable, by digits:

first - type    1, Real (floating point)  
                  2, Integer  
                  4, Boolean  
                  5, Complex

second - shape and mode    0, Scalar  
                             1, Scalar function  
                             2, Vector  
                             3, Vector function  
                             4, Matrix  
                             5, Matrix function  
                             6, Array  
                             7, Array function

third - 0, not a parameter  
          1, non-scalar parameter  
          2, scalar parameter

After the internal symbol table a list of programs used is given. If a program is in the library, its name is prefixed by 'GENIE'. If a name is used, this is given. If a number is used, this is given.

For each definition set a table of external names is printed in which only the names and descriptive information (as above) are of interest.

Compilation options. See PLACER-TRANSLATE.

## RUNNING GENIE PROGRAMS

The usual procedure is to run Genie programs with SPIREL so that all library routines are immediately available.

The initial version of a program should contain liberal output of intermediate quantities. These may be conditioned on sense light settings or edited out once the program is running.

Initial runs should be made with SL14 off so that printer monitoring is provided for all SPIREL operations.

Debugging may be facilitated by a SPIREL dump of the positive portion of the Symbol Table-Value Table. This will show all named external items in the system being run, the values of scalars, and the codewords for non-scalars.

A SPIREL dump of a private program will show values of internal variables.

Arithmetic error tracing may help to locate mathematical problems.

All instructions generated by the compiler may be traced, but this is not a recommended procedure.

## CODING EXAMPLES

### • Least Squares

This program computes the coefficients of a polynomial of specified degree which best fits the input data in the least squares sense. The basic method is described in "An Introduction to Numerical Mathematics", Stiefel, E.L., 1963, page 51. The only difference here being the introduction of weighting factors to the data and the use throughout of matrix algebra.

Lines 6 to 13:

Internal integers are declared and then stored into; the number of rows of XDATA and the length of COEFS (the number of coefficients is compared.

Lines 14 to 45:

The size of XDATA is expanded and is filled with the appropriate powers of X.

Lines 46 to 55:

The normal matrix is computed taking the weights into account.

Lines 56 to 67:

The coefficients, theoretical polynomial values, residues, sum of the squares, and the covariance matrix are computed.

<u>Lines</u>	<u>Comments</u>
4	Some of the parameters, the non-scalars, are declared.
6	Notice lower case alphabetic print output for characters beyond 'f'.
14-40	This AP2 code constructs control words for SPIREL to act upon; notice the labelled instruction at line 35.
41-45	Double or nested looping.
47	A matrix transpose is done here.
56-57	Non-scalar multiplications.

<u>Lines</u>	<u>Comments</u>
60	Solution of a system of equations.
63	Line is labelled but not referred to.
67	Use of matrix exponentiation to compute inverse.



```

1  DEFINE
2  MATRIX XDATA, SIGMA
3  VECTOR YDATA, COFFS, YCALC, RESID, WHTS
4  PFIT(XDATA, YDATA, WHTS, COFFS, YCALC, RESID, SQSUM, SIGMA) = SEQ
5  INTEGER .N, .M, .P, .J, .I
6  .N = LENGTH(YDATA)
7  .M = LENGTH(COFFS)
8  .H = ROW(XDATA)
9  .J = .M - .H
10 CC = 0
11 IF .J = 0
12   7   BAJ      XDATA, U+B1
13       LRS      27
14       CLA      ++1150
15       LRS      12
16       CLA      .J, U+B2
17       LRS      15, R+T7
18       TSR      ++126
19       SPF      *END+1
20       CLA      B1, U+B4
21       CRL      15, R+B2
22   -B2   ADD      +34+B2+1, U+B2
23   P3    BNA      B4+1, U+R
24       CRL      15
25       LRL      12
26       CLA      ++0120
27       LRS      12
28   R     CRR      15, U+T6
29   P3    RPA      T5, B2+1
30   T6    TSR      ++126, U+T7
31       SPF      *END+1
32   P2    IF(NZE)TRA +LJOP, B3+1
33   FOR   .J = 1, 1, .N
34   FOR   .I = .H+1, 1, .M
35   XDATA .I, .J = XDATA .I-1, .J * XDATA .H, .J
36   REPEAT
37   REPEAT
38   WATE  EXECUTE VSPACE(RESID, .N)
39   SIGMA = TRAN(XDATA)
40   FOR .I = 1, 1, .N
41   RESID .I = WHTS .I * YDATA .I
42   FOR .J = 1, 1, .M
43   SIGMA .I, .J = SIGMA .I, .J * WHTS .I
44   REPEAT
45   REPEAT
46   WORK  COFFS = XDATA * RESID
47   SIGMA = XDATA * SIGMA
48   COFFS = SOLN(SIGMA, COFFS)

```

4/06/66 14.18

PAGE 2

	YCALC = TRAN(XDATA) * COEFS	61
	RESID = YDATA - YCALC	62
SUM	SQSUM = 0.0	63
	FOR .I = 1, 1, .N	64
	SQSUM = SQSUM + WCHTS .I + RESID .I <sup>2</sup>	65
	REPEAT	66
	SIGMA = SIGMA <sup>-1</sup> * (SQSUM / (.N - .M))	67
END		70
	DEFINE	71
LEAVE		72
		73

```

+BEGIN PROGRAM SEQUENCE
LOOP PROGRAM SEQUENCE
+FOR1 PROGRAM SEQUENCE
+FOR2 PROGRAM SEQUENCE
+RPT2 PROGRAM SEQUENCE
+RPT1 PROGRAM SEQUENCE
WATE PROGRAM SEQUENCE
+FOR3 PROGRAM SEQUENCE
+FOR4 PROGRAM SEQUENCE
+RPT4 PROGRAM SEQUENCE
+RPT3 PROGRAM SEQUENCE
WORK PROGRAM SEQUENCE
SUMM PROGRAM SEQUENCE
+FOR5 PROGRAM SEQUENCE
+RPT5 PROGRAM SEQUENCE
END PROGRAM SEQUENCE

```

PFIT . =

```

+BEGIN 1 10 01000 02 4400 00136
LOOP 51 43 21601 62 0000 00006
+FOR1 55 20 20001 00 4001 00236 .J
+FOR2 61 20 10000 00 0001 00231 .H
+RPT2 103 20 10401 00 0001 00211 .I
+RPT1 105 20 10401 00 0001 00206 .J
WATE 107 01 21702 26 0200 00005 RESID
+FOR3 124 20 20001 00 4001 00170 .I
+FOR4 137 20 20001 00 4001 00154 .J
+RPT4 155 20 10401 00 0001 00136 .J
+RPT3 157 20 10401 00 0001 00135 .I
WORK 161 01 21700 41 0200 00000 XDATA
SUMM 245 00 20001 00 4600 00006 SQSUM
+FOR5 246 20 20001 00 4001 00046 .I
+RPT5 262 20 10401 00 0001 00032 .I
END 306 01 01000 00 4400 00137
307 01 40006 00 4000 00000
310 07 01000 00 4200 00000

```

REFERENCE WORDS...

```

SMMPY 77770 625454577040000000
MPOWE 77771 545756664440000000
MSUB 77772 546264412540000000
SOLN 77773 625453552540000000
MMPY 77774 545457702540000000
TRAN 77775 636140552540000000
VSPAC 77776 656257474240000000
LENGT 77777 534455446340000000

```

INTERNAL STORAGE..

.N	311	0
.M	312	0
.H	313	0
.J	314	0
.I	315	0
*TW47	316	6200000000000000

# PARAMETERS AT PF +

XDATA	0
YDATA	1
WGHTS	2
COEFS	3
YCALC	4
RESID	5
SQSUM	6
SIGMA	7

# SUBROUTINES REFERENCED

GENIF...	SMMPY	137
GENIF...	MPOWE	
GENIF...	MSUP	
GENIF...	SOLN	
GENIF...	MMPY	
GENIE...	TRAN	135
GENIF...	VSPAC	
GENIF...	LENGT	126
		136

4/20/66 14.27

PAGE 1

PFIT		ORG		BACK-TRANSLATION	1
		REM			2
L77770		REF		*SMMPY	3
L77771		REF		*MPOWE	4
L77772		REF		*MSUB	5
L77773		REF		*SOLN	6
L77774		REF		*MMPY	7
L77775		REF		*TRAN	10
L77776		REF		*VSPAC	11
L77777		REF		*LENGT	12
L1	-Z	TRA		*136, U-R	13
L2		SPF		B4-22	14
	FF	RWT		L307	15
		CLA		PF+1, U-T7	16
		TSR		*L77777	17
		SPF		*L307	20
		STO		L311	21
		CLA		PF+3, U-T7	22
		TSR		*L77777	23
		SPF		*L307	24
		STO		L312	25
		CLA		PF, U-T7	26
		TSR		*L77777	27
		SPF		*L307	30
		STO		L313	31
		21740		L312	32
		SUB		L313	33
		STO		L314	34
		21740		L314	35
		IF(ZER)SKP		a7	36
		TRA		L30	37
		CLA		aL107	40
		NOP		Z, U-CC	41
L30	7	BAU		PF, U-R1	42
		LRS		33	43
		CLA		a1150	44
		LRS		14	45
		CLA		L314, U-B2	46
		LRS		17, R-T7	47
		TSR		*126	50
		SPF		*L307	51
		CLA		B1, U-B4	52
		CRL		17, R-B3	53
	-B2	ADD		a33+B4+1, U-B3	54
	P3	00204		B4+1, U-R	55
		CRL		17	56
		LRL		14	57
		CLA		a120	60
		LRS		14	61
	P	CRR		17, U-T6	62
L51	P3	RPA		T6, B2-1	63
	T6	TSR		*126, U-T7	64
		SPF		*L307	65
	P2	IF(NZF)TRA		L31, B2+1	66
	J	STO		L314	67
L56		CLA		L311	70
		IF(POC)SKP		L314	71
		TRA		L107	72
					73

4/20/66 14.27

PAGE 2

	I	ADD	L313	74
		STO	L315	75
L63		CLA	L312	76
		IF(POS)SKP	L315	77
		TRA	L105	100
		CLA	L314, U→T6	101
	I	BUS+2	L315, B6+1	102
	T6	NOP	Z, U→B2	103
		21740	B4-1, U→B1	104
		21740	*PF, B4=1	105
		NOP	Z, U→T4	106
	T6	NOP	Z, U→B2	107
		21740	L313, U→B1	110
		21740	*PF	111
		10620	T4, U→P	112
	T6	NOP	Z, U→B2	113
		CLA	L315, U→B1	114
	R	STO	*PF	115
	I	FAD→	L315	116
		TRA	L63	117
L105	I	FAD→	L314	120
		TRA	L36	121
L107		CLA+2	PF+5, B6+1	122
	Z	BAU+2	aL311, B6+1	123
		TSR	*L77776	124
		SPF	*L307	125
		CLA	PF, U→T7	126
		TSR	*L77775	127
		SPF	*L307	130
		CLA	PF+7, U→B1	131
	Z	TSR	*L35, U→B2	132
		SPF	*L307	133
	Z	LDR→	10, R→B2	134
	R	STO	B1	135
	B1	RWT	B2	136
	I	STO	L315	137
L125		CLA	L311	140
		IF(POS)SKP	L315	141
		TRA	L161	142
		21740	L315, U→B1	143
		21740	*PF+2, U→T4	144
		21740	L315, U→B1	145
		21740	*PF+1	146
		10620	T4, U→P	147
		CLA	L315, U→B1	150
	R	STO	*PF+5	151
	I	STO	L314	152
L145		CLA	L312	153
		IF(POS)SKP	L314	154
		TRA	L157	155
		CLA	L315, U→T6	156
		CLA	L314, U→B2	157
		21740	T6, U→B1	160
		21740	*PF+7, U→T4	161
		21740	T6, U→B1	162
		21740	*PF+2	163
		10620	T4, U→P	164
		CLA	L314, U→B2	165
	T6	NOP	Z, U→B1	166

4/20/66 14.27

PAGE 3

	F	STO	*PF+7	167
	I	FAD→	L314	170
		TRA	L140	171
L157	I	FAD→	L315	172
		TRA	L125	173
L161		CLA	PF,U→H1	174
		CLA	PF+5,U→B2	175
		TSR	*L77774	176
		SPF	*L307	177
		CLA	PF+3,U→B1	200
	Z	TSR	*135,U→B2	201
		SPF	*L307	202
	Z	LDR→	10,R→H2	203
	F	STO	B1	204
	F1	RWT	B2	205
		CLA	PF,U→B1	206
		CLA	PF+7,U→B2	207
		TSR	*L77774	210
		SPF	*L307	211
		CLA	PF+7,U→B1	212
	Z	TSR	*135,U→B2	213
		SPF	*L307	214
	Z	LDR→	10,R→H2	215
	F	STO	B1	216
	F1	RWT	B2	217
		CLA+2	PF+7,H6+1	220
		CLA+2	PF+3,H6+1	221
		TSR	*L77773	222
		SPF	*L307	223
		CLA	PF+3,U→B1	224
	Z	TSR	*135,U→B2	225
		SPF	*L307	226
	Z	LDR→	10,R→H2	227
	F	STO	B1	230
	F1	RWT	B2	231
		CLA	PF,U→T7	232
		TSR	*L77775	233
		SPF	*L307	234
		CLA	PF+3,U→B2	235
	Z	TSR	*L77774,U→B1	236
		SPF	*L307	237
		CLA	PF+4,U→B1	240
	Z	TSR	*135,U→B2	241
		SPF	*L307	242
	Z	LDR→	10,R→H2	243
	F	STO	B1	244
	F1	RWT	B2	245
		CLA	PF+1,U→B1	246
		CLA	PF+4,U→B2	247
		TSR	*L77772	250
		SPF	*L307	251
		CLA	PF+5,U→B1	252
	Z	TSR	*135,U→B2	253
		SPF	*L307	254
	Z	LDR→	10,R→H2	255
	F	STO	B1	256
	F1	RWT	B2	257
	Z	STO	*PF+6	260
	I	STO	L315	261

4/20/64 14.27

PAGE 4

L247		CLA	L311	262
		IF(POS)SKP	L315	263
		TRA	L264	264
		CLA	L315,U→T6	265
		21740	T6,U→H1	266
		21740	*PF+5	267
		FMP	U,J→T4	270
		21740	T6,U→H1	271
		21740	*PF+2	272
		FAD	T4	273
		FAD→	*PF+6	274
	I	FAD→	L315	275
		TRA	L247	276
L264		21740	L311	277
		SUB	L312	300
		53100	-J	301
		FMP	L316	302
		VDF	*PF+6,U→T4	303
		LDR	-31	304
		CLA	PF+7	305
		TSR	*L77771	306
		SPF	*L307	307
		SB1	10	310
	T4	TSR	*L77770	311
		SPF	*L307	312
		CLA	PF+7,U→B1	313
	Z	TSR	*135,U→B2	314
		SPF	*L307	315
	Z	LDR→	10,R→B2	316
	F	STD	B1	317
	R1	RWT	B2	320
		TRA	*137	321
L307		SB6	Z	322
	T7	TRA	PF	323
L311		OCT	000000000000000000	324
L312		OCT	000000000000000000	325
L313		OCT	000000000000000000	326
L314		OCT	000000000000000000	327
L315		OCT	000000000000000000	330
L316		OCT	042000000000000000	331
		END		332
				333
				334



• Numerical Integration

This example is adapted from Schwarz (An Introduction to ALGOL 60. Comm ACM 5:82-95 (1962)). It concerns the numerical integration of a differential equation of second order with given initial values. Schwarz chose the method of Adams' extrapolation, which consists of the following formulae:

$$y(x+h) = y(x) + hy'(x) + h^2 \left[ \frac{1}{2}y''(x) + \frac{1}{6}\nabla y''(x) + \frac{1}{8}\nabla^2 y''(x) + \dots \right]$$

$$y'(x+h) = y'(x) + h \left[ y''(x) + \frac{1}{2}\nabla y''(x) + \frac{5}{12}\nabla^2 y''(x) + \dots \right]$$

where the  $\nabla^k y''(x)$  are the backward differences of  $y''$  at the point  $x$  and for the interval  $h$ . In contrast to other proposals, he starts the integration by an iterative process (lines 62 to 74) which uses the same formulae as the forward integration (lines 76 to 123).

The example consists of three separate programs: EXAMPLE3, a control program to handle input and output and execute the integration program; F, the function being integrated; and ADAMS, the numerical integration routine. EXAMPLE3 activates STEX and initiates output with a page restore and heading print, then goes into a loop in which it reads four input data from paper tape, performs the integration, prints the input and results, and returns to read more data. ADAMS receives X0,Y0,Z0, and XE as input (with the dummy names XX,YY,ZZ, and EE). M, H, and the final results X,Y, and ZED are external to both EXAMPLE3 and ADAMS.

The integration is based on the following procedure: The leading row of backward differences (which are unknown at the beginning) is first filled out with zeroes (line 52). With this leading row we integrate M steps ahead with the formulae of Adams (line 64), since R in the loop named ADMINT means the number of steps to be integrated. After this we may build up a new difference table from the  $M^{\text{th}}$  row backwards by keeping the  $M^{\text{th}}$  difference con-

stant (lines 67-73). In this way we obtain a new leading row of backward differences, with which we again integrate M steps forward. This is repeated until the  $M^{\text{th}}$  difference of two successive runs are nearly equal (lines 65-66 and 74; note that WE is the  $M^{\text{th}}$  difference of the preceding run). As soon as BETA is FALSE, we start integrating ahead a sufficient number of steps to reach XE (lines 76 to 123).

<u>Lines</u>	<u>Comments</u>
13-16	An AP2 sequence is used to initialize output and activate STEX.
61-63	Note use of the power point in arithmetic expressions.
17,23,24, 27-34	Input and results are printed with SCRIBE. The arguments in the EXECUTE command correspond in number and order to the dummy fields in the FORMATS.
11-12, 42-46	The REM may be followed by either a tab (lines 42-46) or a carriage return (lines 11-12). The same is true of FORMAT and BCD.
52-60	Extra spaces are ignored.
37	This line illustrates both the definition of a function in a single line and the use of an auxilliary equation to evaluate a common sub-expression.
51	Execution of VSPACE leaves zeroes in the vector for which space is taken. This initializes W for the first pass through the loop.
61,65,66, 61,63	If a name occurs for the first time on the lefthand side of an equation, its type is inferred from the righthand side. Thus, DECIDE and BETA are inferred to be Boolean in lines 61 and 63; R, J, and V are inferred as integers in lines 61,65, and 66.
63	BETA is evaluated as TRUE if $10^{-7} <  W_M - WE $ ; otherwise BETA is evaluated as FALSE.

<u>Lines</u>	<u>Comments</u>
72,76,105 122	These are all conditional equations. Lines 76 and 105 illustrate arithmetic conditionals; lines 72 and 122 illustrate Boolean conditionals.
22&41 102&37	The values to be used at each execution of a function are passed to the function as an ordered argument list, with the arguments corresponding in number and type to the parameters in the definition of the program.
123	The vectors for which space was taken at the beginning of ADAMS are freed at the end.

```

1
2  DEFINE
3  SCALARS XO,YO,ZO,XF,H,X,Y,ZED
4  INTEGER M
5  VECTORS B,C,W
6  FUNCTIONS F,ADAMS
7
8  EXAMPLE3(7),=SEQ
9  REM
10 THIS IS THE DRIVER PROGRAM. IT CONTROLS INPUT, INTEGRATION, AND OUTPUT
11
12     PAG      +Z
13     SLN      +00002
14     LT7      +00000 3120 0000 00135
15     TSR      *+126
16
17 EXECUTE SCRIBE(HEADER)
18 M=6, H=0.01
19
20 LOOP DATA XO,YO,ZO,XE
21 EXECUTE ADAMS(XO,YO,ZO,XE)
22 EXECUTE SCRIBE(M,H,XO,YO,ZO,IN)
23 EXECUTE SCRIBE(X,Y,ZF,OUT)
24     SPA      +7
25
26 CC=-LOOP
27
28 HEADER FORMAT
29 M      H      YO      YO      ZO
30
31 IN      FORMAT
32 d      d, dddd -dddd, dddd -dddd, dddd -dddd, dddd
33 OUT     FORMAT
34 -dddd, dddd -dddd, dddd -dddd, dddd
35
36 END
37
38 F(X,Y,Z)=ZZ(SIN(TMP)+COS(TMP)+YY2+2), TMP=XX+YY+ZZ
39
40 ADAMS(XX,YY,ZZ,EE),=SEQ
41 REM XX,YY,ZZ ARE THE INITIAL VALUES FOR X,Y,Z PRIME
42 REM M IS THE ORDER OF THE METHOD (≤6)
43 REM EE IS THE END OF THE INTEGRATION
44 REM H IS THE INTEGRATION STEP
45 REM W0 IS THE SECOND DERIVATIVE, Wk THE KTH BACK DIFF.
46
47 EXECUTE VSPACE(B,7)
48 EXECUTE VSPACE(C,7)
49 EXECUTE VSPACE(W,M+1)
50
51 P1=1., C1=0.5
52
53 P2=0.5, C2=1./6.
54
55 P3=5./12., C3=1./8.
56
57 P4=3./8., C4=19./180.
58
59 P5=251./720., C5=3./32.
60

```

4/19/66 13.24

PAGE 2

	R <sub>6</sub> =95./288.,	C <sub>6</sub> =963./10080.	57
	R <sub>7</sub> =19087./60480.,	C <sub>7</sub> =275./3456.	60
LOOP	WE=1*10, R=M, DECIDE=TRUE		61
RLINT	CC=+ADMINT		62
	BETA=1*-7<1W <sub>M</sub> -WE1		63
	WE=W <sub>M</sub>		64
	FOR J=M,-1,1		65
	FOR V=2,1,M		66
	W <sub>V</sub> =W <sub>V</sub> -W <sub>V+1</sub>		67
	REPEAT		70
	REPEAT		71
	CC=+LOOP .If BETA		72
ADMINT	R=FIX((1X-XX)/H), DECIDE=FALSE		73
	X=XX, Y=YY, ZED=7Z		74
	FOR J=1,1,R+1		75
	CC=+L14 .If J=1		76
	FOR V=M,-1,1		77
	W <sub>V+1</sub> =W <sub>V</sub>		100
	REPEAT		101
L14	W <sub>1</sub> =F(X,Y,ZED)		102
	REM F IS THE FUNCTION DEFINING THE		103
	DIFFERENTIAL EQUATION		104
	CC=+NSHIFT .If J=C		105
	FOR V=2,1,M+1		106
	W <sub>V+1</sub> =W <sub>V</sub>		107
NSHIFT	REPEAT		110
	P=Z, Q=Z		111
	FOR V=1,1,M+1		112
	P=P+B <sub>V</sub> W <sub>V</sub>		113
	Q=Q+C <sub>V</sub> W <sub>V</sub>		114
	REPEAT		115
	X=X+H		116
	Y=Y+H(ZED+Q H)		117
	ZED=ZED+P H		120
	REPEAT		121
	CC=+RLINT .If DECIDE		122
	FRASE B,C,W		123
END			124
LEAVE	DEFINE		125
			126
			127

EXAMP START NEW PROGRAM

4/19/66 13.33

+BGIN PROGRAM SEQUENCE  
 LOOP PROGRAM SEQUENCE  
 HEADE PROGRAM SEQUENCE  
 IN PROGRAM SEQUENCE  
 OUT PROGRAM SEQUENCE  
 END PROGRAM SEQUENCE

EXAMP . =

+BGIN 1 47 21641 00 0001 00101 END

THIS IS THE DRIVER PROGRAM. IT CONTROLS INPUT, INTEGRATION, AND OUTPUT

LOOP 15 01 40001 00 4000 00004  
 HEADE 51 00 00000 00 0000 00007  
 IN 61 00 00500 00 0000 00010  
 OUT 72 00 00300 00 0000 00010  
 END 103 01 01000 00 4000 00000

REFERENCE WORDS...

ED 77764 714443252500000000  
 Y 77765 702525252500000000  
 X 77766 672525252500000000  
 ADAMS 77767 404340546240000000  
 XE 77770 674425252500000000  
 ZO 77771 710025252500000000  
 YO 77772 700025252500000000  
 XO 77773 670025252500000000  
 +INQU 77774 755055556644000000  
 H 77775 472525252500000000  
 M 77776 542525252500000000  
 SCRIB 77777 624261504140000000

INTERNAL STORAGE..

+NUMB 104 3120000000135  
 +NUMB 105 770024363605075341

SUBROUTINES REFERENCED

ADAMS  
 GENIE... +INQU  
 GENIE... SCRIB

F START NEW PROGRAM

4/19/66 13.33

+BGIN PROGRAM SEQUENCE  
END PROGRAM SEQUENCE

F . =

+BGIN	1	10 01000 02 4400 00136
END	25	01 01000 00 4400 00137
	26	01 40004 00 4000 00000
	27	07 01000 00 4200 00000
REFERENCE WORDS...		
COS	77776	425662232540000000
SIN	77777	#250552525400000000
INTERNAL STORAGE..		
TMP	30	0
+NUMB	31	101200000000000000

PARAMETERS AT PF +

XX	0
YY	1
ZZ	2

SUBROUTINES REFERENCED

GENIE...	COS	137
GENIE...	SIN	
		136

ADAMS START NEW PROGRAM

4/19/66 13.33

+BGIN PROGRAM SEQUENCE  
 LOOP PROGRAM SEQUENCE  
 RLINT PROGRAM SEQUENCE  
 +FOR1 PROGRAM SEQUENCE  
 +FOR2 PROGRAM SEQUENCE  
 +RPT2 PROGRAM SEQUENCE  
 +RPT1 PROGRAM SEQUENCE  
 ADMIN PROGRAM SEQUENCE  
 +FOR3 PROGRAM SEQUENCE  
 +FOR4 PROGRAM SEQUENCE  
 +RPT4 PROGRAM SEQUENCE  
 L14 PROGRAM SEQUENCE  
 +FOR5 PROGRAM SEQUENCE  
 +RPT5 PROGRAM SEQUENCE  
 NSHIF PROGRAM SEQUENCE  
 +FOR6 PROGRAM SEQUENCE  
 +RPT6 PROGRAM SEQUENCE  
 +RPT3 PROGRAM SEQUENCE  
 END PROGRAM SEQUENCE

ADAMS . =

+BGIN 1 10 01000 02 4400 00136

XX,YY,ZZ ARE THE INITIAL VALUES FOR X,Y,Y PRIME

M IS THE ORDER OF THE METHOD (≤4)

EE IS THE END OF THE INTEGRATION

H IS THE INTEGRATION STEP

W+0 IS THE SECOND DERIVATIVE, W+K THE KTH BACK DIFF.

LOOP	121	01	21700	40	4001	00053	ADMIN
RLINT	122	01	21740	41	0401	77650	M
+FOR1	134	01	21700	00	0401	77636	M
+FOR2	141	01	21700	00	4000	00002	
+RPT2	155	20	10401	00	0001	00230	V
+RPT1	157	30	10401	00	0001	00225	J
ADMIN	175	01	21700	00	0600	00000	XX
+FOR3	202	20	20001	00	4001	00201	J
+FOR4	215	01	21700	00	0401	77555	M
+RPT4	230	30	10401	00	0001	00155	V
L14	232	00	20102	26	4401	77534	X

F IS THE FUNCTION DEFINING THE  
 DIFFERENTIAL EQUATION

+FOR5	244	01	21700	00	4000	00002	
+RPT5	263	20	10401	00	0001	00122	V
NSHIF	265	00	20001	00	4001	00122	P
+FOR6	267	20	20001	00	4001	00116	V
+RPT6	313	20	10401	00	0001	00072	V
+RPT3	330	20	10401	00	0001	00054	J



END	347	01 01000 00 4400 00137
	350	01 40004 00 4000 00000
	351	07 01000 00 4200 00000

# REFERENCE WORDS...

F	77764	452525252540000000
ZED	77765	714443252500000000
Y	77766	702525252500000000
X	77767	672525253500000000
FIX	77770	455067252540000000
H	77771	472525252500000000
XE	77772	674425252500000000
M	77773	542525252500000000
W	77774	662525252540000000
C	77775	422525252540000000
VSPAC	77776	656257414240000000
B	77777	412525252540000000

# INTERNAL STORAGE..

+P1	352	0
+ONEF	353	100100000000000000
+NUMB	354	772000000000000000
+NUMB	355	100600000000000000
+NUMB	356	100500000000000000
+NUMB	357	101400000000000000
+NUMB	360	101000000000000000
+NUMB	361	100300000000000000
+NUMB	362	100300000000000000
+NUMB	363	120400000000000000
+NUMB	364	130300000000000000
+NUMB	365	200264000000000000
+NUMB	366	104000000000000000
+NUMB	367	110700000000000000
+NUMB	370	200110000000000000
+NUMB	371	200327600000000000
+NUMB	372	204730000000000000
+NUMB	373	211243000000000000
+NUMB	374	235420000000000000
+NUMB	375	200104600000000000
+NUMB	376	201540000000000000
+NUMB	377	50022500574400000
WE	400	0
R	401	0
DECID	402	0
+NUMB	403	750015327745152745
BETA	404	0
J	405	0
V	406	0
+P2	407	0
P	410	0
Q	411	0

# PARAMETERS AT PF +

XX	0
YY	1
ZZ	2
EE	3

SUBROUTINES REFERENCED

137  
135

GENIF... F  
GENIE... FIX  
VSPAC

136

• Complex Matrix Inverse

This program inverts a square matrix whose elements are complex numbers. The method used is essentially in-place Gaussian reduction as described in "An Introduction to Numerical Mathematics", Stiefel, E.L., 1963, page 3. Each successive pivot element has the largest modulus of all the remaining choices. This insures the least possible error in the resulting inverse. If the modulus of any pivot element is too small, the matrix is numerically singular, an error message is printed.

The  $n \times n$  complex matrix is stored as  $2 n \times n$  real matrices with the primary codewords in two successive memory locations. Throughout the program, subscripting and arithmetic are performed on the complex variables with the same Genie code that has heretofore been used for real variables.

Lines 11 to 13:

Working storage defined.

Line 14:

Complex matrix B is copied into A. Thus, B will be preserved after its inverse is computed.

Lines 20 to 27:

The largest remaining pivot element is found and stored in GMOD and indices stored in GG and HH.

Lines 30 to 64:

If the chosen pivot is large enough, the exchange algorithm is applied to A.

Lines 66 to 103:

Since pivot elements were not in general along the diagonal, the rows and columns of the inverse are rearranged depending on the contents of ROWW and COLL.

Line 104:

The inverse is stored in RESULT, where all results of implicit

functions are stored. Working storage is freed.

<u>Lines</u>	<u>Comments</u>
3-4	Double declarations of integer vectors and complex matrices.
10	Use of ROW function implicitly in expression.
11-12	ROWW and COLL are declared real and, thus, are created by VSPACE, the real vector space program called in the program.
13	NEW is declared complex. Thus, even though MSPACE is called by the user, CMSPACE (complex matrix space) will be the program executed.
22	Subscripting of a complex variable; use of MOD function implicitly in an expression.
24	Two equations on one line, as many as fifteen permitted.
30	Specification of a constant using power point, '*'. '←' is printed for '#'.
45	Use of '+/' to create complex variable out of two real variables.
51	Labelled REPEAT command. This is not the same as labelling the corresponding FOR statement.
60	Compound conditional transfer.
106	Unconditional transfer to labelled location.
107	Use of SCRIBE to print error message.
115-117	Terminating LEAVE statement is followed by two carriage returns.

9/23/66 13.02

PAGE 1

	DEFINE	1
	COMPLEX MATRIX B,A,NEW	2
	INTEGER VECTOR ROWW,COLL	3
	INVERT(B),=SEQ	4
	INTEGER GG,HH,L,C,D,E	5
	COMPLEX G	6
	L=ROW(B)	7
	EXECUTE VSPACE(ROWW,L)	10
	EXECUTE VSPACE(COLL,L)	11
	EXECUTE MSPACE(NEW,L,L)	12
	A=B	13
	FOR C=1,1,L	14
	GMOD=0.	15
	GG=C,HH=C	16
	FOR D=C,1,L	17
	FOR E=C,1,L	20
	KMOD=MOD(A <sub>D,E</sub> )	21
	CC=+MORE .If KMOD≤GMOD	22
	GG=D,HH=E	23
	GMOD=KMOD	24
MORE	REPEAT	25
	REPEAT	26
	CC=+BYE .If GMOD≤1.0*-12	27
	COLL <sub>C</sub> =GG	30
	ROWW <sub>C</sub> =HH	31
	FOR D=1,1,L	32
	G=A <sub>C,D</sub>	33
	A <sub>C,D</sub> =A <sub>GG,D</sub>	34
	A <sub>GG,D</sub> =G	35
	REPEAT	36
	FOR D=1,1,L	37
	G=A <sub>D,C</sub>	40
	A <sub>D,C</sub> =A <sub>D,HH</sub>	41
	A <sub>D,HH</sub> =G	42
	REPEAT	43
	NEW <sub>C,C</sub> =(1.+/-0.)/A <sub>C,C</sub>	44
	FOR D=1,1,L	45
	CC=+SOME .If D=C	46
	NEW <sub>D,C</sub> =A <sub>D,C</sub> /A <sub>C,C</sub>	47
SOME	REPEAT	50
	FOR D=1,1,L	51
GENIE	CC=+TOME .If D=C	52
September, 1966	NEW <sub>C,D</sub> =-(A <sub>C,D</sub> /A <sub>C,C</sub> )	53
		54

9/23/66 13.02

PAGE 2

TOME	REPEAT	55
	FOR D=1,1,L	56
	FOR E=1,1,L	57
	CC=VOME .If D=C .O.R E=C	60
	NEW D,E=A D,E+(NEW C,F)(A D,C)	61
VOME	REPEAT	62
	REPEAT	63
	A=NEW	64
	REPEAT	65
	FOR C=L,-1,1	66
	HH=COLL C	67
	GG=ROWW C	70
	FOR D=1,1,L	71
	G=A D,C	72
	A D,C=A D,HH	73
	A D,HH=G	74
	REPEAT	75
	FOR D=1,1,L	76
	G=A C,D	77
	A C,D=A GG,D	100
	A GG,D=G	101
	REPEAT	102
	REPEAT	103
	RESULT=A	104
	ERASE COLL,ROWW,NEW,A	105
	CC=END	106
BYE	EXECUTE SCRIBE(MESS)	107
	CC=END	110
MESS	FORMAT	111
NO INVERSE	DUE TO SINGULARITY	112
END		113
	DEFINE	114
LEAVE		115
		116

INVER START NEW PROGRAM

9/23/66 13.02

+BGIN PROGRAM SEQUENCE  
 +FOR1 PROGRAM SEQUENCE  
 +FOR2 PROGRAM SEQUENCE  
 +FOR3 PROGRAM SEQUENCE  
 MORE PROGRAM SEQUENCE  
 +RPT2 PROGRAM SEQUENCE  
 +FOR4 PROGRAM SEQUENCE  
 +RPT4 PROGRAM SEQUENCE  
 +FOR5 PROGRAM SEQUENCE  
 +RPT5 PROGRAM SEQUENCE  
 +FOR6 PROGRAM SEQUENCE  
 SOME PROGRAM SEQUENCE  
 +FOR7 PROGRAM SEQUENCE  
 TOME PROGRAM SEQUENCE  
 +FOR8 PROGRAM SEQUENCE  
 +FOR9 PROGRAM SEQUENCE  
 VOME PROGRAM SEQUENCE  
 +RPT8 PROGRAM SEQUENCE  
 +RPT1 PROGRAM SEQUENCE  
 +FORa PROGRAM SEQUENCE  
 +FORb PROGRAM SEQUENCE  
 +RPTb PROGRAM SEQUENCE  
 +FORc PROGRAM SEQUENCE  
 +RPTc PROGRAM SEQUENCE  
 +RPTa PROGRAM SEQUENCE  
 BYE PROGRAM SEQUENCE  
 MESS PROGRAM SEQUENCE  
 END PROGRAM SEQUENCE

INVER :=

+BGIN	1	10 01000 02 4400 00136	
+FOR1	47	20 20001 00 4001 00511	C
+FOR2	60	01 21700 00 0001 00500	C
+FOR3	65	01 21700 00 0001 00473	C
MORE	115	20 10401 00 0001 00450	E
+RPT2	117	20 10401 00 0001 00445	D
+FOR4	134	20 20001 00 4001 00430	D
+RPT4	162	20 10401 00 0001 00402	D
+FOR5	164	20 20001 00 4001 00400	D
+RPT5	212	20 10401 00 0001 00352	D
+FOR6	232	20 20001 00 4001 00331	D
SOME	265	20 10401 00 0001 00277	D
+FOR7	267	20 20001 00 4001 00275	D
TOME	321	20 10401 00 0001 00243	D
+FOR8	323	20 20001 00 4001 00241	D
+FOR9	327	20 20001 00 4001 00236	E

GENIE  
 September, 1966

VOME	37F	20 10401 00 0001 00170	E
+RPT8	377	20 10401 00 0001 00165	D
+RPT1	420	20 10401 00 0001 00140	C
+FORa	422	01 21700 00 0001 00133	L
+FORb	43F	20 20001 00 4001 00127	D
+RPTb	462	20 10401 00 0001 00101	D
+FORc	46F	20 20001 00 4001 00077	D
+RPTc	512	20 10401 00 0001 00051	D
+RPTa	51F	30 10401 00 0001 00043	C
BYE	541	01 21702 26 4001 00004	MESS
MESS	546	00 00000 00 0000 00004	
END	553	01 01000 00 4400 00137	
	554	01 40006 00 4000 00000	
	555	07 01000 00 4200 00000	

# REFERENCE WORDS...

SCRIB	77755	424261504140000000
CADD	77754	424043432540000000
CMFY	77757	425457702540000000
CDIV	77760	424250652540000000
CMPLX	77761	425457536700000000
+++++	77762	757575757500000000
MOD	77762	545443252540000000
CSTAR	77764	426263476140000000
+++++	77765	757575757540000000
A	77766	402525252540000000
+++++	77767	757575757540000000
CMCPY	77770	425442577040000000
CMSPA	77771	425462574040000000
NEW	77772	554466252540000000
+++++	77772	757575757540000000
COLL	77774	425453532540000000
VSPAC	77775	456257404240000000
ROWW	77776	615466625400000000
CROW	77777	426156662540000000

# INTERNAL STORAGE..

L	554	0
+P1	557	0
+P2	560	0
C	561	0
GMOD	562	0
GG	562	0
HH	564	0
D	565	0
E	566	0
+Q1	567	0
+++++	570	0
KMOD	571	0
+NUMB	572	730010627463004557
G	572	0
+++++	574	0
+ONEF	575	100100000000000000

GENIE

September, 1966

PARAMETERS AT PF +

B	0
+++++	1



SUBROUTINES REFERENCED

GENIE...	SCRIB	137
GENIE...	CADD	
GENIE...	CMFY	
GENIE...	CDIV	
GENIE...	MOD	
GENIE...	CMCPY	135
GENIE...	CMSPA	
GENIE...	VSPAC	
GENIE...	CROW	136

## CODING CONVENTIONS

This section discusses details of compiler generated code. It is intended for those who are particularly interested and for those who wish to code in a lower level language while maintaining compatibility with compiled programs. This material is not essential to the understanding of the Genie language and should not be read before attempting to write some programs for the compiler and gaining some familiarity with the Rice Computer, the assembly language, and the SPIREL system.

- Programs initialization and termination

The 'SEQ' or 'RSEQ' causes the compiler to generate a sequence of orders which initializes the program being compiled. The first of these orders is labelled '←BGIN', and the orders are collectively called the "←BGIN code sequence". For each 'SEQ' or 'RSEQ' there is an 'END', and an "END code sequence" corresponds to each ←BGIN code sequence. The forms of these code sequences depend on whether 'SEQ' or 'RSEQ' is used, the number of parameters (p) listed for the program and, in some cases, the types of the parameters. Each complex parameter is counted as two parameters, the real part followed by the imaginary part.

An 'SEQ' causes generation of a non-recursive program; an 'RSEQ' causes generation of a recursive program. These two types of code are distinguished functionally by the location of internal variables for the program. Constants are always stored within the program. Private storage is inside a non-recursive program and on the B6-list, addressed relative to PF, for a recursive program. Genie-generated recursive code will not alter itself while running, and a recursive program may use itself -- provided AP2 code in the program also obeys conventions necessary for recursion. The use of a program by itself is clear in a case where program A uses program A; if program A uses B which uses C which uses A, then again program A is using itself.

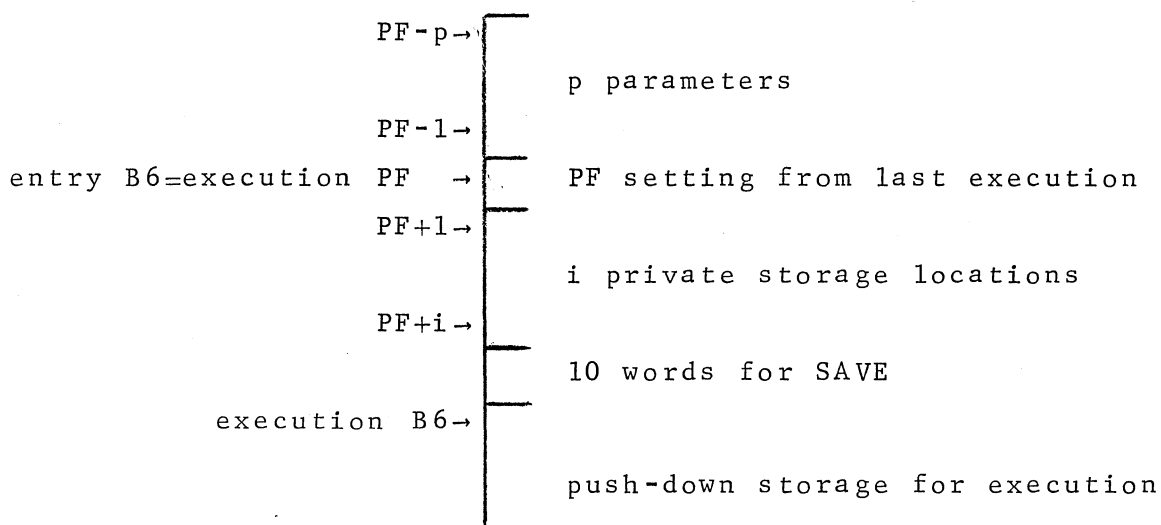
For a non-recursive program -- one begun with 'SEQ' or a one statement function ... A single fast parameter in the definition of a program is a special case which causes only PF to be saved and assumes no parameter addressing in Genie language within the program. Otherwise, fast register names should not be used as parameters in a program definition, and the following discussion applies. A single parameter enters a program in T7, the value of a scalar or \* codeword address for a non-scalar. Immediately a scalar with name P in T7 is stored at internal location 'P'; a non-scalar parameter is stored on the B6-list. All fast registers

are saved; if there are parameters on the B6-list ( $p > 1$  or  $p = 1$  and a non-scalar parameter) PF is set to point to the first parameter. In this case (PF) is stored in the address portion of 'END+1' and must be maintained with this value throughout the program for the purpose of addressing parameters. The END code sequence restores the fast registers, sets B6 to free the storage occupied by any parameters on the B6-list, fetches (T7) for implicit execution, and exits to the PF setting on entry. The specific code sequences are as follows:

$p = 1$ fast	←BGIN	PF	RWT : : :	END
	END		TRA	Z
$p = 1$ scalar	←BGIN	-Z T7	TRA STO : : :	*+136, U→R P
	END	T7	TRA TRA	*+137 PF
$p = 1$ non-scalar	←BGIN	T7 -Z PF	STO TRA SPF RWT : : :	B6, B6+1 *+136, U→R B6-11 END+1
	END	T7	TRA SB6 TRA	*+137 (Z) PF
$p > 1$	←BGIN	-Z PF	TRA SPF RWT : : :	*+136, U→R B6-p-10 END+1
	END	T7	TRA SB6 TRA	*+137 (Z) PF

For a recursive program -- one begun with 'RSEQ' ... A single parameter enters a program in T7, the value of a scalar or \* codeword address for a non-scalar. Multiple parameters enter on the B6-list at B6-p, ..., B6-1, address for a scalar and \* codeword address for a non-scalar. A single non-scalar parameter is stored on the B6-list. In all cases the PF setting for the last execution of the program is picked up from 'END+1' and stored just

beyond the parameters on the B6-list. This B6 value is stored in 'END+1' for the PF setting of the current execution. B6 is advanced over  $i$  private storage locations for the program. A full save is done. Then PF is set for execution -- with  $p$  parameters at PF-p,...,PF-1 and  $i$  private storage locations at PF+1,...,PF+i. A single scalar parameter named P is stored at private storage location 'P'. In the case of a single fast or a single scalar parameter, the program is considered to have no parameters. B6-list utilization by a recursive program is illustrated by:



The END code sequence restores all fast registers, restores the PF setting for the last execution at END+1, backs B6 up by  $p+i+1$  to free all B6-list list used in execution, fetches (T7) for implicit execution, and exits to the PF setting on entry. The specific code sequences are as follows:

## CODING CONVENTIONS

5

single scalar (p=0)	←BGIN	B6	CLA,WTG+2	END+1
		B6	RWT	END+1,B6+1
		-Z	ADD	a - *END+3,U→B6
			TRA	*+136,U→R
			SPF	*END+1
		T7	STO	P
single non-scalar (p=1)	←BGIN	T7	STO	B6,B6+1
			CLA,WTG+2	END+1
		B6	RWT	END+1,B6+1
		B6	ADD	a - *END+3,U→B6
		-Z	TRA	*+136,U→R
			SPF	*END+1
single fast (p=0) and multiple (p>1)	←BGIN		CLA,WTG+2	END+1
		B6	RWT	END+1,B6+1
		B6	ADD	a - *END+3,U→B6
		-Z	TRA	*+136,U→R
			SPF	*END+1
all cases	END		TRA	*+137
			CLA,WTG	Z
			STO,WTG	END+1,B6-1
		PF	AB6	[-i],U→R
		T7	AB6	[-p],R→CC

- Result for implicit execution

A program which is single valued may be executed implicitly; that is, it may be mentioned within the formula on the righthand side of an equation in Genie language. A non-complex scalar result must be in U upon exit from the program, a complex scalar result in the complex accumulator named CMPLX, a non-complex non-scalar result in the non-scalar accumulator whose codeword is by definition at location +10 during execution. The name 'RESULT' is interpreted by the compiler as T7 for a non-complex scalar, as CMPLX for a complex scalar, as codeword address +10 for a non-complex non-scalar and as CSTAR for a complex non-scalar. 'RESULT' may appear only on the lefthand side of an equation and must be defined in the last command executed before 'END' on all dynamic paths to 'END'. The 'END' code sequence fetches (T7) to U as it exits so that a non-complex scalar result is indeed in U upon return to the program causing the implicit execution.

- Addressing of variables

With respect to any given program every variable is in one of three categories: internal, external, parameter. All internal variables are scalar. For a non-recursive program the values of all internal variables are stored within the program. For a recursive program internal variables are of two types: constants are stored within the program; others are stored in private storage on the B6-list, the  $i^{\text{th}}$  private storage word being addressed at  $(\text{PF})+i$  after program initialization. External variables may be scalar or non-scalar, the address or \* codeword address respectively being stored in a cross reference word within the program, the value or codeword respectively being stored in the Value Table (\*+122) during execution. In the general case, reference words for parameters are stored on the B6-list. For a non-recursive program the  $p^{\text{th}}$  parameter is addressed at  $(\text{PF})+p-1$  after program initialization. For a recursive program the  $p^{\text{th}}$  parameter is addressed at  $(\text{PF})-p$  after program initialization. Parameters of a program during execution are indeed internal or external with respect to some dynamically higher level program, but this does not affect addressing in the program where they are parameters. The following charts summarize addressing conventions for variables.



## CODING CONVENTIONS

8

For a non-recursive program --

<u>variable</u>	<u>representation</u>	<u>data address</u>	<u>codeword address</u>	<u>value</u>	<u>element</u>
internal scalar	value in program at IS	#IS	-----	(IS)	-----
external scalar	address in program at ES	(ES)	-----	*ES	-----
external non-scalar	* codeword address in program at ENS	-----	address in (ENS)	-----	*ENS
scalar parameter	address at PF+p-1	(PF+p-1)	-----	*PF+p-1	-----
non-scalar parameter	* codeword address at PF+p-1	-----	address in (PF+p-1)	-----	*PF+p-1

For a recursive program --

<u>variable</u>	<u>representation</u>	<u>data address</u>	<u>codeword address</u>	<u>value</u>	<u>element</u>
internal constant	value in program at IC	#IC	-----	(IC)	-----
internal storage	value on B6-list at PF+i	#PF+i	-----	(PF+i)	-----
external scalar	address in program at ES	(ES)	-----	*ES	-----
external non-scalar	* codeword address in program at ENS	-----	address in (ENS)	-----	*ENS
scalar parameter	address on B6-list at PF-p	(PF-p)	-----	*PF-p	-----
non-scalar parameter	* codeword address on B6-list at PF-p	-----	address in (PF-p)	-----	*PF-p

• B6-list, working storage

The SPIREL system uses the block with codeword address +112 as a working storage area. The conventions associated with this storage are that B6 points to the next available location on the list [hence, the term "B6-list"] and that the storage is used in a linear "last-in-first-out" or "push-down" fashion. Genie generated code uses the B6-list for temporary storage of intermediate quantities within the calculation of an arithmetic formula, always storing at (B6), incrementing (B6) after the store, retrieving from (B6)-1, and decrementing (B6) after retrieval. The B6-list is also used for storage of parameters before entering a program; the program then decrements (B6) over the parameters before return since the storage occupied by parameters is no longer in use. For a recursive program, a private storage area is established on the B6-list and freed prior to exit. The SAVE (\*+136) and UNSAVE (\*+137) programs and other SPIREL routines use the B6-list for temporary dynamic push-down storage.

Using the B6-list for temporary storage, the following sequence shows storage of A, B, C and later retrieval of C, B, A with proper maintenance of (B6) as a pointer to the B6-list:

```

:
CLA+2    A, B6+1
:
CLA+2    B, B6+1
:
CLA+2    C, B6+1
:
: calculation perhaps involving
: use of B6-list with balance of
: stores and retrivals, so that
: final (B6) = initial (B6)
CLA      B6-1, B6-1
STO      C
:
CLA      B6-1, B6-1
STO      B
:
CLA      B6-1, B6-1
STO      A
:

```

- Parameter set-up for program execution

Execution of a program with a single non-complex scalar parameter SP is preceded by code which accomplishes (SP)→T7. In the case of a single non-scalar parameter NSP, the code accomplishes \*NSP→T7. For more than one parameter, representations are stored sequentially on the B6-list; if the  $k^{\text{th}}$  parameter is a scalar SP, then SP→B6, B6+1; if the  $k^{\text{th}}$  parameter is a non-scalar NSP, then \*NSP→B6, B6+1. A complex parameter is treated as two parameters, the real part followed by the imaginary part. If one of a group of parameters is given by a number or an expression, then the quantity must be given a name before the proper parameter representation can be stored on the B6-list. For such purpose the names '←P1', '←P2', etc. for non-complex quantities are generated by the compiler. The quantity is stored at ←Pn for a scalar or \*←Pn for a non-scalar is stored on the B6-list. A non-scalar at ←Pn is freed upon return from the program for which it was stored; then all ←Pn used are available for re-use. Complex quantities are stored as pairs named '←Q1', '←Q2', etc., then each part is treated like a non-complex parameter.

The execution of program PROG is accomplished by TSR \*PROG where PROG is a cross-reference word for PROG within the program doing the execution; the codeword for PROG is in the Value Table (\*+122). Thus, PROG is an external variable with respect to the program which executes it.

- Representation of Complex Variables

A complex variable is always on the first level of addressing represented by a pair of words in consecutive memory locations, the real part followed by the imaginary part. The name of a complex variable is attached to the first word of the pair, the real part; the second word of the pair has the name "ditto", printed '←←←←'. The Cartesian form is used, and both parts are real floating point.

Genie generates internal storage for the complex scalar A as

```

A          real part of A
←←←←      imaginary part of A

```

Genie generates cross-reference words for the external complex variable A as

```

A          name 'A' in hexads/* if non-scalar/
          VT address for A
←←←←      name "ditto" in hexads/* if non-scalar/
          VT address for A's ditto

```

Then while running the corresponding ST-VT configuration is

ST	VT
A ...	real part of A - value if scalar, codeword if non-scalar
←←←← ...	imaginary part of A - value if scalar, codeword if non-scalar

Genie constructs two argument words on the B6-list for each complex argument A. The first addresses the real part of A; the second addresses the imaginary part of A.

- Subscription

In the Genie language any variable may be subscripted by from one to five indices separated by commas. The indices are assumed by the compiler to be integers: explicit numbers, simple names, or arithmetic expressions of any complexity. The indices are loaded successively into B1, B2, ..., B5 by the following procedure which allows subscripts to themselves be subscripted:

- 1) scan n indices from left to right, computing those which are not numbers or simple names, and storing those computed (except the last) on the B6-list;
- 2) scan from right to left storing (U), quantity from B6-list, named quantity, or explicit number into  $B_i$  for  $i=n, n-1, \dots, 1$ .

In the sense of SPIREL, a subscripted variable is called an "array". In particular, a one-dimensional array of data is called a "vector" and is indexed by B1, and a two-dimensional array of data is called a "matrix" and is indexed by B1 and B2 in that order. But in fact an array may be of as many as five dimensions and may contain either data or programs, and its elements may be addressed in the Genie language. The indices may take on negative values if the storage configuration is correspondingly established.

- Operations on standard forms of non-scalars

In order to perform an operation between a scalar and a vector or matrix, to combine two vectors or matrices, or to store a vector or matrix the non-scalar itself must be addressed in the code. Although completely general forms of non-scalars may be created and manipulated in the SPIREL context and may have their elements addressed in the Genie language, operations on full vectors and matrices are defined only for arrays of standard form in order that execution time is not spent in handling the most general case. The standard form of non-scalars is entirely sufficient in a vast majority of applications. The definition is as follows:

standard form of one dimensional array, vector

- 1) loaded with STEX active
- 2) indexed by B1
- 3) initial index = 1

standard form of two dimensional array, matrix

- 1) loaded with STEX active
- 2) indexed by B1 for row specification and B2 for column specification
- 3) initial row index = 1, initial column index = 1

A standard complex non-scalar is a pair of standard non-scalars, as described. Codewords must be adjacent, real then imaginary; a name adheres to the real part, and the imaginary part is named "ditto" (←←←←).

Arithmetic operations involving standard non-scalars parallels scalar arithmetic quite closely. By convention, codeword +10 is used as the non-complex non-scalar accumulator, commonly called 'U\*'; the complex non-scalar accumulator is named CSTAR. The programs used for performing operations on non-scalars recognize a null codeword address for a non-scalar operand to mean that the operand is the accumulator. The creation of a new U\* or CSTAR causes the storage previously addressed by that "name" to be freed. If a non-scalar in U\* or CSTAR needs to be temporarily saved, this is done on the B6-list; that is, a word or pair of words on the

B6-list are taken as codewords for the storage addressed and the accumulator codewords are cleared. Note that this storage also involves adjustment of the STEX back-references to address the new codewords.

The code sequence generated by the compiler for non-complex non-scalar storage  $A \rightarrow B$  is as follows:

	CLA	A, U→B2	}	copy A→U* only if $A \nsubseteq U^*$
Z	TSR	*MCOPY, U→B1		
≠	SPF	*END+1		
	CLA	B, U→B1	}	free storage addressed as B only if $B \nsubseteq U^*$ and not on B6-list
Z	TSR	*+135, U→B2		
≠	SPF	*END+1		
Z	LDR→	+10, R→B2	}	clear U* codeword store new codeword for B update back-reference
R	STO	B1		
B1	RPA, WTG	B2		
				if $B \nsubseteq U^*$

The code sequence generated by the compiler for complex non-scalar storage  $A \rightarrow B$  is as follows:

	CLA	A, U→B2	}	copy A→CSTAR only if $A \nsubseteq CSTAR$
Z	TSR	*CMCOPY, U→B1		
≠	SPF	*END+1		
	CLA, DBL	B, R→B1	}	free storage addressed as B only if $B \nsubseteq CSTAR$ and not on B6-list
Z	TSR	*+135, U→B2		
	NOP	Z, B1-1		
Z	TSR	*+135, U→B2	}	store new codewords for B
	CLA	CSTAR, U→PF		
	CLA, DBL	PF, U→B2		
	STO, DBL	B1	}	update back-references
B1	RPA	B2, R→B2		
	NOP	Z, B1+1		
B1	RPA	B2, R→Z	}	clear CSTAR codewords
Z	STO, DBL	PF		
≠	SPF	*ENDP1		
				if $B \nsubseteq CSTAR$

≠(PF) reset only if program is recursive or is using (PF) for reference to parameters.

• Assignment of type and shape to variables

In the Genie language each variable has a shape: scalar, vector, or matrix. The shape of a variable may be explicitly specified as non-scalar by a declaration: VECTOR for vector, MATRIX for matrix. Each scalar, vector, matrix, and function (result) has a type: integer, real floating point, complex, or Boolean. The type of a variable may be explicitly specified in a declaration: INTEGER for integer, REAL or SCALAR for real floating point, COMPLEX for complex, and BOOLEAN for Boolean. The standard shape/type is scalar/floating point unless otherwise specified in an INFER declaration. If the first appearance of a variable name is not in a declaration, its type is implicitly specified by the following rules:

- 1) If a variable name first appears on the right side of an equation, the variable is assigned the standard shape/type.
- 2) If a variable name first appears on the lefthand side of an equation, the variable is assigned the shape/type of the expression on the righthand side.

In a compilation a variable will not have its type changed once it is assigned. An equation which has lefthand and righthand sides of different types will cause the compiler to comment on the equating of unlike types; code will be generated to perform a store appropriate to the quantity on the righthand side, but the type of the quantity on the lefthand side will be unaffected.



- Arithmetic combination of variables of different types

In arithmetic expressions Boolean and integer variables may be combined only in exponentiation, Boolean scalar variable to an integer scalar power. Boolean and floating point variables may not be combined.

Integer and real floating point scalars and non-scalars may be combined in any mathematically meaningful way. In all cases except exponentiation of a floating point scalar by a numerically specified integer  $\leq 7$ , the integer must be floated before the combination takes place. In all cases the result of the combination is floating point. If a numerically defined integer scalar is floated, the floating point equivalent is generated at compilation time and is referenced in the generated code for the combination. Otherwise, the floating of an integer scalar A is Accomplished by the following generated code:

```

-LDU      -A
FMP      ←TW47

```

where '←TW47' refers to the constant  $2^{47}$  which will be stored within the program. The floating of an integer vector or matrix is accomplished by use of the Genie SPIREL program MFLT.

Integers and real floating point scalars and non-scalars may be combined with complex scalars and non-scalars in any mathematically meaningful way. In all cases except exponentiation of a complex scalar by an integer or floating point scalar the non-complex quantity is made complex before the combination takes place. A floating point quantity is made complex with real part equal the floating point quantity and zero imaginary part; an integer quantity is floated then made complex as a floating point quantity.

- Boolean variables and operations

A Boolean variable may take on the value 'TRUE' or 'FALSE', these being represented in the computer by full length quantities

TRUE = +007777777777777777

FALSE = +007777777777777776

The binary operations between Boolean variables to yield a Boolean value cause code to be generated as follows:

or, A+B, true if either A or B is true

CLA A

ORU B

and, A×B, true if both A and B are true

CLA A

ORU B

symmetric difference, A-B, true if A and B have different values

CLA A

SYD B

ORU #77776

symmetric sum, A/B, true if A and B have the same value

CLA -A

SYD B

The only meaningful unary operation on a Boolean variable is complementation, not A, true if A is false

-I ORU -A

The machine register sense lights (SL) is a collection of 15 bits, any one of which may be individually meaningful and may be in an on or off (1 or 0) state at any time. The variable SL is Boolean and exponentiation to an integer power is defined

$A^B$ , true if bit B of A is on (1) where the bits of A are numbered from 1 to 15, from left to right

CLA	A	}	
LUR	15-B		if B is a number
ORU	#+77776	}	
CLA	B		
BUS	#15,U→R		if B is
CLA	A		a name
			or
LUR	*R		an expression
ORU	#+77776	}	

Although the Boolean exponential notation is particularly meaningful for the lights, it may be applied to any Boolean variable. Thus, a Boolean variable A which does not itself have a value of TRUE or FALSE may be a collection of 15 bits (the rightmost in a machine word)  $A^1, A^2, \dots, A^{15}$  each with a value of TRUE or FALSE.

- Loop coding

In the Genie language a loop is begun by the command

FOR iteration parameter = initial, increment, final and  
ended by the command

REPEAT

If there are not labels on these commands, the  $k^{\text{th}}$  loop will have the labels ' $\leftarrow\text{FOR}k$ ' and ' $\leftarrow\text{RPT}k$ ' associated with it. The generalized code generated for loop control is as follows:

$\leftarrow\text{FOR}k$	compute initial		
	initial $\rightarrow$ iteration parameter		
	compute increment		
	store increment		
	compute final		
	store final		
$[\leftarrow\text{FOR}k+m]$	LT7	final	
Z	IF(POS)SKP	increment	
T7	IF(POS)SKP	iteration parameter, CC+1	
T7	IF(NEG)SKP	iteration parameter	
	TRA	$\leftarrow\text{RPT}k+n$	
	:		
	orders of loop		
	:		
$\leftarrow\text{RPT}k$	CLA	increment	
	FAD $\rightarrow$	iteration parameter	
	TRA	$\leftarrow\text{FOR}k+m$	
	:		
$[\leftarrow\text{RPT}k+n]$	:		

Seldom is the full generalized code necessary, and the following notes pertain to condensations which are provided in various specific cases.

- (A) The increment and the final value are computed and stored only if they are given by expressions, that is, not simple variable names or explicit numbers.

- (B) The final value will be stored in the address field of the order if it is given by an explicit integer.
- (C) If the increment is given by an explicit integer, it will not be tested for being positive or negative and only the appropriate comparison of iteration parameter to final value will be generated.
- (D) If the iteration parameter is a long fast register F, the  $\leftarrow$ RPTk code sequence will be
- |                   |   |     |                              |
|-------------------|---|-----|------------------------------|
| $\leftarrow$ RPTk | F | FAD | increment, U $\rightarrow$ F |
|                   |   | TRA | $\leftarrow$ FORk+m          |

If the iteration parameter is an index register Bi and the increment is an explicit integer +1 or -1, the  $\leftarrow$ RPTk code sequence will be

$\leftarrow$ RPTk		TRA	$\leftarrow$ FORk+m, Bi $\pm$ 1
-------------------	--	-----	---------------------------------

• Use of fast registers in Genie generated code

Fast registers may be used in the Genie language and in assembly language coding to be used in a Genie context if there is no conflict with usage generated by the compiler:

T7 is always subject to use for special purpose temporary storage.

T7 is used for storage of a single parameter when a function is executed implicitly or explicitly.

T4, T5, T6 are subject to use in any arithmetic command for scalar temporary storage and for storage of scalars mentioned two or more times in one equation if these fast register names are not mentioned explicitly in the command.

B1 is used when loading parameters onto the B6-list if a name  $\leftarrow P_n$  is used.

B1, B2, B3, B4, B5 are used for subscripts in addressing elements of arrays. The first k are used to address an element of an array of k dimensions.

B1 and B2 are used in complex scalar arithmetic.

B1, B2, and PF may be used in operations on vectors and matrices.

B1 is used in input-output commands to specify to the program  $\leftarrow INOUT$  the operation to be performed.

B1 is used in raising an integer or a real floating point scalar to an integer power  $\leq 7$ .

B6 always addresses the push-down B6-list which is used for temporary storage of scalars and non-scalars, for multiple parameter storage, and for private storage of a recursive program.

PF is used within a non-recursive program to address its parameters if there are more than one or if there is only one but that is a non-scalar. The appropriate value for (PF) is, in such cases, stored in the address portion of END+1 so that resetting is easily accomplished by

SPF                      \*END+1

PF is used within every recursive program to address parameters and private storage locations. The appropriate value of (PF) is stored in the address portion of END+1 so that resetting is easily accomplished by

SPF                      \*END+1

- Rearrangement of arithmetic formulae for efficient evaluation

The compiler has the ability to rearrange the terms in addition (or subtraction) and multiplication (or division) strings. Constant terms are shifted to the left in the formula. Terms which are themselves expressions, rather than simple variable names or numbers, are shifted to the left to save temporary stores that would be required were such complex terms to appear to the right in a string. The ordering of the complex terms is determined by the number of temporary stores required to evaluate each; the complex term requiring the most temporary stores will be shifted farthest to the left.

If the order of evaluation within a formula is of importance, this rearrangement may be avoided by defining each complex term in a separate equation, thereby giving each a name. Then the original formula will involve only simple variable names, and rearrangement will not take place.